

**University of Nevada, Las Vegas Computer Science 456/656 Fall 2021**

**Assignment 3. Due 11:30 AM Monday September 20, 2021**

**Name:**-----

You are permitted to work in groups, get help from others, read books, and use the internet.

Do not hand this homework in. It will not be graded, but I will go over it in class on September 20.

1. Starting on page 138 of the sixth edition of your textbook or on page 137 of the fifth edition, work problems 9(b), 9(e), 12(c), and 15.

2. I will give the proof by contradiction that  $\sqrt{2}$  is irrational in class. Prove by contradiction that  $\sqrt{3}$  is irrational.

3. You have probably seen the easy (high school) proof that the sum of the first  $n$  positive integers is  $\frac{n(n+1)}{2}$ . That formula was made famous by Johann Carl Friedrich Gauss, who came up with the proof at age 10 to work a class assignment in his head in less than a minute while the other pupils were busy for an hour. He was not the first to find the formula, which was well-known at the time. The teacher knew it, but didn't tell the kids. There is also a proof by induction, which I will give in class.

Use induction to prove that the sum of the squares of the first  $n$  positive integers is  $\frac{n(n+1)(2n+1)}{6}$ . What is the sum of the cubes of the first  $n$  positive integers? Look it up.

The context-free grammar  $G_1$  given below generates an algebraic language  $L$ , where the strings of  $L$  are C++ expressions with variables and operators. There are only three C++ variables,  $x$ ,  $y$ , and  $z$ , and there are only four operations, addition, subtraction, negation, and multiplication. The start symbol is  $E$  (for “expression”) instead of the usual  $S$ , and the alphabet of  $L$  is  $\{x, y, z, +, -, *, (, )\}$

(a) Prove that  $G_1$  is ambiguous by giving two different left-most derivations of the string  $x + y * z$  and the corresponding derivation trees.

1.  $E \rightarrow E + E$
2.  $E \rightarrow E - E$
3.  $E \rightarrow E * E$
4.  $E \rightarrow -E$
5.  $E \rightarrow (E)$
6.  $E \rightarrow x$
7.  $E \rightarrow y$
8.  $E \rightarrow z$

(b) When I input a certain string  $w \in L$  into my top-down parser, it outputs 352674513687, the encoding of a left-most derivation of  $w$ . Find  $w$ .

(c) The grammar  $G_1$  does not “respect” the C++ precedence of operators. (What does that mean?) Give an unambiguous grammar  $G_2$  for  $L$  which respects C++ precedence rules. Use three grammar variables:  $E$  for “expression,”  $T$  for “term,” and  $F$  for “factor.” I will help get you started:

1.  $E \rightarrow E + T$
2.  $E \rightarrow E - T$
3.  $E \rightarrow T$
4.  $T \rightarrow T * F$
5.  $T \rightarrow$  (Missing rhs.)
6.  $\rightarrow -F$  (Missing lhs.)
7.  $F \rightarrow ( \ )$  (What goes inside the parentheses?)
8.  $F \rightarrow x$
9.  $F \rightarrow y$
10.  $F \rightarrow z$

4. Fill in the blanks.

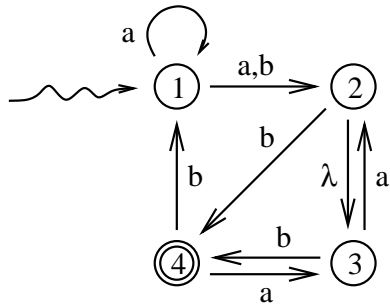
- (a) If  $M_1$  is an NFA with  $n$  states, then  $M_1$  is equivalent to a DFA  $M_2$  with no more than \_\_\_\_\_ states.
- (b) The \_\_\_\_\_ algorithm determines whether a given string of length  $n$  is generated by a given context-free grammar in  $O(n^3)$  time.
- (c) A derivation tree of a string  $w$  generated by context-free grammar  $G$  can be labeled by writing, next to each variable symbol in the tree, the number of a production of  $G$ . The encoding of a left-most derivation of  $w$  is obtained by visiting the internal nodes of the corresponding derivation tree in \_\_\_\_\_. (The answer is either preorder, postorder, inorder, or level order.)

5. State the pumping lemma for regular languages by filling in the blanks.

If  $L$  is a \_\_\_\_\_ language, there exists an integer  $p$ , which we call the \_\_\_\_\_ of  $L$ , such that for every \_\_\_\_\_ such that \_\_\_\_\_, there exist strings \_\_\_\_\_, \_\_\_\_\_ and \_\_\_\_\_ such that the following four statements hold:

- 1. \_\_\_\_\_
- 2. \_\_\_\_\_
- 3. \_\_\_\_\_
- 4. For any \_\_\_\_\_, the string \_\_\_\_\_ is a member of  $L$ .

6. Design a minimal DFA equivalent to the following NFA.



7. Design a PDA which accepts the language  $L$  of all even length palindromes over  $\{a, b\}$ ,

$$S \rightarrow aSa$$

$$S \rightarrow bSb$$

$$S \rightarrow \lambda$$

which is generated by this context-free grammar. You may describe your machine either by writing the transitions or by drawing a figure.

8. Let  $L$  be the language consisting of all binary strings with equally many zeros and ones. For example, 0100011011 and 1001 are in  $L$ , but 01101001100101101 is not. Design a DPDA which accepts  $L$ . You may describe your machine either by writing the transitions or by drawing a figure.

9. Let  $L = \{a^n b^n c^n : n \geq 0\}$ . As you know,  $L$  is not context-free.

(a) Give two context-free languages whose intersection is  $L$ .

(b) Give a context-sensitive grammar for  $L$ .

(c) Give a context-free grammar for the complement of  $L$ . (This is hard. You need several variables.)

10. True or False. T = true, F = false, and O = open, meaning that the answer is not known science at this time. In the questions below,  $\mathcal{P}$  and  $\mathcal{NP}$  denote  $\mathcal{P}$ -TIME and  $\mathcal{NP}$ -TIME, respectively.
- (i) \_\_\_\_\_ Let  $L$  be the language over  $\{a, b, c\}$  consisting of all strings which have more  $a$ 's than  $b$ 's and more  $b$ 's than  $c$ 's. There is some PDA that accepts  $L$ .
  - (ii) \_\_\_\_\_ The language  $\{a^n b^n \mid n \geq 0\}$  is context-free.
  - (iii) \_\_\_\_\_ The language  $\{a^n b^n c^n \mid n \geq 0\}$  is context-free.
  - (iv) \_\_\_\_\_ The language  $\{a^i b^j c^k \mid j = i + k\}$  is context-free.
  - (v) \_\_\_\_\_ The intersection of any three regular languages is regular.
  - (vi) \_\_\_\_\_ The intersection of any regular language with any context-free language is context-free.
  - (vii) \_\_\_\_\_ The intersection of any two context-free languages is context-free.
  - (viii) \_\_\_\_\_ The complement of any context-free language is context-free.
  - (ix) \_\_\_\_\_ If  $L$  is a context-free language over an alphabet with just one symbol, then  $L$  is regular.
  - (x) \_\_\_\_\_ The set of strings that your high school algebra teacher would accept as legitimate expressions is a context-free language.
  - (xi) \_\_\_\_\_ Every language accepted by a non-deterministic machine is accepted by some deterministic machine.
  - (xii) \_\_\_\_\_ The problem of whether a given string is generated by a given context-free grammar is decidable.
  - (xiii) \_\_\_\_\_ If  $G$  is a context-free grammar, the question of whether  $L(G) = \emptyset$  is decidable.
  - (xiv) \_\_\_\_\_ Every language generated by an unambiguous context-free grammar is accepted by some DPDA.
  - (xv) \_\_\_\_\_ The language  $\{a^n b^n c^n d^n \mid n \geq 0\}$  is decidable.
  - (xvi) \_\_\_\_\_
  - (xvii) \_\_\_\_\_ The language  $\{a^n b^n c^n \mid n \geq 0\}$  is in the class  $\mathcal{P}$ -TIME.
  - (xviii) \_\_\_\_\_ There exists a polynomial time algorithm which finds the factors of any positive integer, where the input is given as a binary numeral.
  - (xix) \_\_\_\_\_
  - (xx) \_\_\_\_\_ Every problem that can be mathematically defined has an algorithmic solution.
  - (xxi) \_\_\_\_\_
  - (xxii) \_\_\_\_\_ Every  $\mathcal{NP}$  language is decidable.

- (xxiii) \_\_\_\_\_ The intersection of two  $\mathcal{NP}$  languages must be  $\mathcal{NP}$ .
- (xxiv) \_\_\_\_\_  $\mathcal{P} = \mathcal{NP}$ .
- (xxv) \_\_\_\_\_  $\mathcal{NP} = \mathcal{P}$ -SPACE
- (xxvi) \_\_\_\_\_ Primality, where the input is in binary, is  $\mathcal{P}$ -TIME.
- (xxvii) \_\_\_\_\_ Every context-free language is in  $\mathcal{P}$ .
- (xxviii) \_\_\_\_\_ The problem of whether two given context-free grammars generate the same language is decidable.
- (xxix) \_\_\_\_\_ If  $G$  is a context-free grammar, the question of whether  $L(G) = \Sigma^*$  is decidable, where  $\Sigma$  is the terminal alphabet of  $G$ .
- (xxx) \_\_\_\_\_ The language of all fractions (using base 10 numeration) whose values are less than  $\pi$  is decidable. (A *fraction* is a string. “314/100” is in the language, but “22/7” is not.)
- (xxxi) \_\_\_\_\_ There exists a polynomial time algorithm which finds the factors of any positive integer, where the input is given as a unary (“caveman”) numeral.
- (xxxii) \_\_\_\_\_ If  $L$  is  $\mathcal{NP}$  and also  $\text{co-}\mathcal{NP}$ , then  $L$  must be  $\mathcal{P}$ .
- (xxxiii) \_\_\_\_\_ There exists a mathematical proposition that is true but cannot be proved.
- (xxxiv) \_\_\_\_\_ There is an uncomputable function which grows faster than any computable function.
- (xxxv) \_\_\_\_\_ There exists a machine that runs forever and outputs the decimal digits of  $\pi$  (the well-known ratio of the circumference of a circle to its diameter).
- (xxxvi) \_\_\_\_\_ For every real number  $x$ , there exists a machine that runs forever and outputs the decimal digits of  $x$ .
- (xxxvii) \_\_\_\_\_ There is a polynomial time algorithm which determines whether any two regular expressions are equivalent.