

Instantaneous Description (id)

An **id** is a string, the concatenation of the following three strings.

1. The contents of the tape before the current cell, possibly the empty string.
2. The current state.
3. The contents of the current cell, one symbol.
4. The contents of the tape after the current cell but with trailing blanks deleted; possibly the empty string.

We denote a blank by “ \sqcup .” A state q_i for an integer i , and q_1 is the start state. Each transition are labeled as $a \rightarrow b, M$, where a is the symbol at the current cell, b is symbol written at the current cell, and M is the movement of the read-write head, either L, R or S , for left, right, and stay. If $a = b$, we write $a \rightarrow M$ for short. The input string w is initially in the first $|w|$ cells, and the leftmost cell is the initial current cell. Thus, the initial **id** is $q_1 w$.

Example

Figure 1 illustrates a TM M which decides the language $\{0^{2^k}\}$. The input alphabet is $\{0\}$ and the tape alphabet is $\{0, x, \sqcup\}$. If the input is accepted, M halts in state q_6 , otherwise it halts in state q_7 .

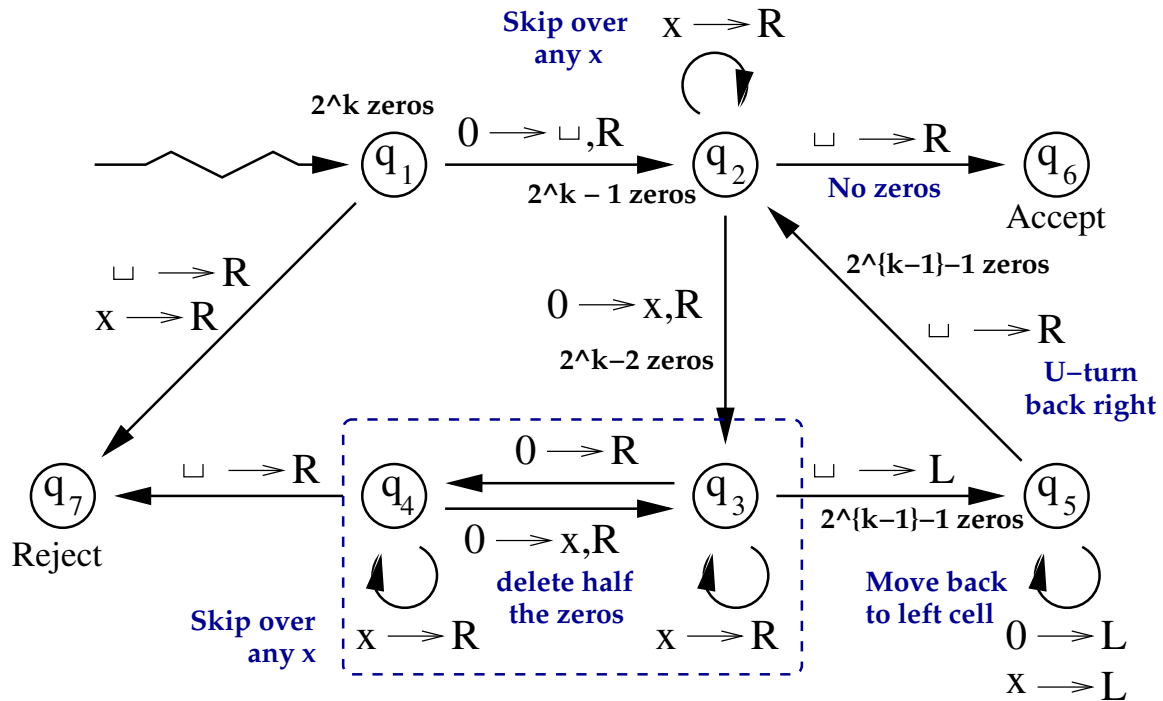


Figure 1: State Diagram of M , a Turing Machine which Decides $\{0^{2^k}\}$

Explanation

At each state during an accepting computation, that is, if the input is 0^{2^k} , the number of zeros in the string is indicated in the diagram. The first zero is removed from q_1 to q_2 . If there is none, the input is rejected. A zero is removed by replacing it by x .

An accepting computation iterates a loop from q_2 back to q_2 . If there are m zeros at q_2 , and m is odd, the back and forth movement between q_3 and q_4 deletes $\frac{m+1}{2}$ zeros. If $m = 2^k$, after k iterations of the loop, the computation reaches q_2 with no zeros on the tape, and the next step is to accept by moving to q_6 . On the other hand, if the initial number of zeros is not a power of 2, eventually there will be no zeros on the tape at q_4 , and the next state is q_7 , the reject state.

For example: if the input is 0^{16} , there will be 15 zeros at q_2 , then later 7 zeros at q_2 , then 3, then 1, then 0, followed by acceptance. On the other hand, if the the input is 0^{12} , there will at first be 11 zeros on the tape at q_2 , then 5 the next time, then 2 the next time, leading to the current cell being blank at q_4 , followed by rejection. We give computations of M with input 0^n for values of n from 1 to 6.

$q_1 0$
□ q_2 □
□□ q_6 □
Accept

$q_1 00$
□ $q_2 0$
□ xq_3 □
□ $q_5 x$
 q_5 □ x
□ $q_2 x$
□ xq_2 □
□ x □ q_6 □
Accept

$q_1 000$
□ $q_2 00$
□ $xq_3 0$
□ $x 0q_4$ □
□ $x 0$ □ q_7 □
Reject

$q_1 0000$
□ $q_2 000$
□ $xq_3 00$
□ $x 0q_4 0$
□ $x 0xq_3$ □
□ $x 0q_5 x$
□ $xq_5 0x$
□ $q_5 x 0x$
 q_5 □ $x 0x$

$\sqcup q_2 x 0 x$
 $\sqcup x q_2 0 x$
 $\sqcup x x q_3 x$
 $\sqcup x x x q_3 \sqcup$
 $\sqcup x x q_5 x$
 $\sqcup x q_5 x x$
 $\sqcup q_5 x x x$
 $q_5 \sqcup x x x$
 $\sqcup q_2 x x x$
 $\sqcup x q_2 x x$
 $\sqcup x x q_2 x$
 $\sqcup x x x q_2 \sqcup$
 $\sqcup x x x \sqcup q_6 \sqcup$
Accept

$q_1 0 0 0 0 0$
 $\sqcup q_2 0 0 0 0$
 $\sqcup x q_3 0 0 0$
 $\sqcup x 0 q_4 0 0$
 $\sqcup x 0 x q_3 0$
 $\sqcup x 0 x 0 q_4 \sqcup$
 $\sqcup x 0 x 0 \sqcup q_7 \sqcup$
Reject

$q_1 0 0 0 0 0 0$
 $\sqcup q_2 0 0 0 0 0$
 $\sqcup x q_3 0 0 0 0$
 $\sqcup x 0 q_4 0 0 0$
 $\sqcup x 0 x q_3 0 0$
 $\sqcup x 0 x 0 q_4 0$
 $\sqcup x 0 x 0 x q_3 \sqcup$
 $\sqcup x 0 x 0 q_5 x$
 $\sqcup x 0 x q_5 0 x$
 $\sqcup x 0 q_5 x 0 x$
 $\sqcup x q_5 0 x 0 x$
 $\sqcup q_5 x 0 x 0 x$
 $q_5 \sqcup x 0 x 0 x$
 $\sqcup q_2 x 0 x 0 x$
 $\sqcup x q_2 0 x 0 x$
 $\sqcup x x q_3 x 0 x$
 $\sqcup x x x q_3 0 x$
 $\sqcup x x x 0 q_4 x$
 $\sqcup x x x 0 x q_4 \sqcup$
 $\sqcup x x x 0 x \sqcup q_7 \sqcup$
Reject