## University of Nevada, Las Vegas Computer Science 456/656 Spring 2019 Assignment 4: Due March 6, 2019

## Name:\_\_\_\_\_

Print this document and staple, along with any extra sheets you want graded. Fill in answers by hand, not by typing or by computer. Hand it to the graduate assistant at the beginning of class on March 6. You are permitted to work in groups, get help from others, read books, and use the internet. But the handwriting on this document must be your own.

Recall that a language L is called *decidable* there is some machine M such that if M is given any string w as input, M will give the output "1" if  $w \in L$ , "0" if  $w \notin L$ . Decidable languages are also called *recursive*.

A function is said to be computable, or recursive, if it is computed by some machine. There are uncomputable functions, such as the Busy Beaver function:

## http://googology.wikia.com/wiki/Busy\_beaver\_function

In fact, there are more uncomputable functions than computable functions.

- 1. True or False. T = true, F = false, and O = open, meaning that the answer is not known to science at this time. In the questions below,  $\mathcal{P}$  and  $\mathcal{NP}$  denote  $\mathcal{P}$ -TIME and  $\mathcal{NP}$ -TIME, respectively.
  - (a) \_\_\_\_\_ Every language generated by an unambiguous context-free grammar is accepted by some DPDA.
  - (b) \_\_\_\_\_ The language  $\{a^n b^n c^n d^n \mid n \ge 0\}$  is recursive.
  - (c) \_\_\_\_\_ Let L be the language over  $\{a, b, c\}$  consisting of all strings which have more a's than b's and more b's than c's. There is some PDA that accepts L.
  - (d) \_\_\_\_\_ The language  $\{a^n b^n c^n \mid n \ge 0\}$  is in the class  $\mathcal{P}$ .
  - (e) \_\_\_\_\_ There exists a polynomial time algorithm which finds the factors of any positive integer, where the input is given as a binary numeral.
  - (f) \_\_\_\_\_ The language  $\{a^n b^n \mid n \ge 0\}$  is context-free.
  - (g) \_\_\_\_\_ The language  $\{a^n b^n c^n \mid n \ge 0\}$  is context-free.
  - (h) \_\_\_\_\_ The language  $\{a^i b^j c^k \mid j = i + k\}$  is context-free.
  - (i) \_\_\_\_\_ The intersection of any regular languages with any context-free language is context-free.
  - (j) \_\_\_\_\_ If L is a context-free language over an alphabet with just one symbol, then L is regular.
  - (k)  $\longrightarrow \mathcal{P} = \mathcal{NP}$ .
  - (l) \_\_\_\_\_ There is a deterministic parser for any context-free grammar.
  - (m) \_\_\_\_\_ The Boolean Circuit Problem is in  $\mathcal{P}$ .

- (n) \_\_\_\_\_ The Boolean Circuit Problem is in  $\mathcal{NC}$ .
- (o) \_\_\_\_\_ The set of strings that your high school algebra teacher would accept as legitimate expressions is a context-free language.
- (p) \_\_\_\_\_ There exists a polynomial time algorithm which finds the factors of any positive integer, where the input is given as a binary numeral.
- (q) \_\_\_\_\_ The language consisting of all strings over  $\{a, b\}$  which have more a's than b's is context-free.
- (r) \_\_\_\_\_ Every context-free language is in  $\mathcal{P}$ .
- (s) \_\_\_\_\_ Every language accepted by a non-deterministic machine is accepted by some deterministic machine.
- (t) \_\_\_\_\_ The problem of whether two given context-free grammars generate the same language is decidable.
- (u) \_\_\_\_\_ The problem of whether a given string is generated by a given context-free grammar is decidable.
- (v) \_\_\_\_\_ Every bounded function is recursive.
- (w) \_\_\_\_\_ There exists a mathematical proposition that can be neither proved nor disproved.
- (x) \_\_\_\_\_ There is a non-recursive function which grows faster than any recursive function.
- (y) \_\_\_\_\_ There exists a machine<sup>1</sup> that runs forever and outputs the string of decimal digits of  $\pi$  (the well-known ratio of the circumference of a circle to its diameter).
- (z) \_\_\_\_\_ For every real number x, there exists a machine that runs forever and outputs the string of decimal digits of x.

<sup>&</sup>lt;sup>1</sup>As always in automata theory, "machine" means abstract machine, a mathematical object whose memory and running time are **not** constrained by the size and lifetime of the known (or unknown) universe, or any other physical laws. If we want to discuss the kind of machine that exists (or could exist) physically, we call it a "physical machine."

- 2. The language  $L = \{a^i b^j c^k : i = j \text{ or } j = k\}$  is generated by the CF grammar G with start state S given below.
  - 1.  $S \rightarrow XC$ 2.  $X \rightarrow \varepsilon$ 3.  $X \rightarrow aXb$ 4.  $C \rightarrow \varepsilon$ 5.  $C \rightarrow cC$ 6.  $S \rightarrow AY$ 7.  $A \rightarrow \varepsilon$ 8.  $A \rightarrow aA$ 9.  $Y \rightarrow \varepsilon$ 10.  $Y \rightarrow bYc$

Design a PDA which makes use the grammar G and accepts L. Walk through the steps of the PDA for the input string *abbcc*.