## Computer Science 456/656 Spring 2020

## Second Examination April 30, 2020

The entire examination is 205 points.

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The exam is take-home, open book, open notes, open internet. You must finish by midnight of April 30 . Scan and email the completed examination paper to your TA, Pradip Marahajan. The email must have an April 30 time stamp.

1. True or False. [5 points each] $\mathrm{T}=$ true, $\mathrm{F}=$ false, and $\mathrm{O}=$ open, meaning that the answer is not known to science at this time.
i -------- Every subset of a regular language is regular.
ii $\qquad$ EXP-TIME $\subseteq$ EXP-SPACE.
iii $\qquad$ There exists a context-sensitive language which is $\mathcal{P}$-SPACE complete.
iv $\qquad$ Every finite language is regular.
v $\qquad$ The language $\left\{a^{i} b^{j} a^{j} b^{i}: i, j \geq 0\right\}$ is context-free.
vi $\qquad$ Any languge generated by an unrestricted grammar is recursively enumerable.
vii $\qquad$ Every polynomial time language is context-free.
viii $\qquad$ If $L$ is in $\mathcal{P}$-space, there is a reduction of $L$ to the regular expression equivalence problem.
ix $\qquad$ The union of two undecidable languages is always undecidable.
x $\qquad$ The union of two recursively enumerable languages is always recursively enumerable.
xi $\qquad$ The union of two co-RE languages is always co-RE. Hint: Think!
xii -------- $\mathcal{N C}=c o-\mathcal{N C}$. Hint: Think!
xiii $\qquad$ The set of all regular expressions for regular languages over the alphabet $\{a, b\}$ is a contextfree language.
xiv $\qquad$ Various websites, such as https://www.youtube.com/watch? v=bQnjbDHefgc give solutions to various instances of RUSH HOUR. If there is a solution to a particular instance of RUSH HOUR, that solution can always be explained in polynomial time.
xv --------The factoring problem for an integer written in binary is both $\mathcal{N P}$ and co- $\mathcal{N} \mathcal{P}$.

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xvi $\qquad$ If someone somewhere on the Earth publishes a correct proof that the partition problem is in $\mathcal{P}$-TIME, then it will be known that $\mathcal{P}=\mathcal{N} \mathcal{P}$.
xvii $\qquad$ If someone somewhere on the Earth publishes a correct proof that the factoring problem for binary numerals is in $\mathcal{P}$-Time, then it will be known that $\mathcal{P}=\mathcal{N} \mathcal{P}$.
xviii $\qquad$ $\mathcal{N C}=\mathcal{P}$-SPACE
xix $\qquad$ co- $\mathcal{N P} \subseteq \mathcal{P}$-space.
2. Fill in the blanks. [ 10 points each blank.]
(a) If $L \subseteq \Sigma^{*}$ is $\mathcal{N P}$ time, there is a constant $k$ and a deterministic machine $V$ such that, for string $w \in \Sigma^{*}$, we have $w \in L$ if and only if there is a string $c$, called a $\qquad$ for $w$, such that $|c| \leq|w|^{k}$ and $V$ accepts the string $c w$ within $|w|^{k}$ steps.
(b) The practicality of the RSA one-way encryption system depends on the assumption (which has not been verified) that the $\qquad$ problem cannot be solved in polynomial time.
(c) $\mathcal{N C}$ is the class of languages which can be decided in $\qquad$ time using polynomially many processors.
3. [20 points] Every context-free language has a minimum pumping length. For example, the minimum pumping length of $\left\{a^{n} b^{n}: n \geq 0\right\}$ is 2 .

The language $L=\left\{a^{n} b^{m} c d e^{n}: n, m \geq 0\right\}$ is context-free.
(a) Find the minimum pumping length of $L$. Call it $p$. $\qquad$
(b) For every string $w \in L$ of length at least $p$, there are strings $u, v, x, y, z$ such that $w=u v x y z$ and three other conditions are satisfied. Find the strings $u, v, x, y, z$ if $w=a a b c d e e$.

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4. [20 points] Give a polynomial time reduction of 3-SAT to the independent set problem.

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5. [20 points] Prove that the halting problem is undecidable.

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6. [20 points] Prove that the language $L=\left\{a^{n} b^{m} c^{m} d^{n}: n, m \geq 0\right\}$ is context-free by giving a context-free grammar for $L$.
