

Using the Lock-down Browser

You will have to answer each question using just typing Here are examples.

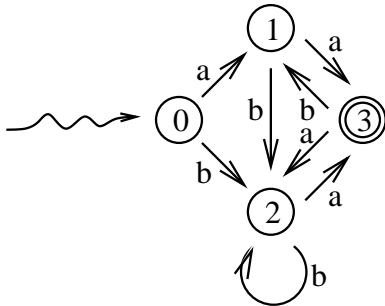
I. Design an NFA which accepts the language of all strings over a, b which have three consecutive a 's.

Your answer:

States are 0, 1, 2, 3 where 3 is the only final state.

$\delta(0, a) = 0, 1$
 $\delta(0, b) = 0$
 $\delta(1, a) = 2$
 $\delta(1, b) = \text{emptyset}$
 $\delta(2, a) = 3$
 $\delta(2, b) = \text{emptyset}$
 $\delta(3, a) = 3$
 $\delta(3, b) = 3$

II. Give a minimal DFA equivalent to the following DFA:



Your answer: States 1 and 2 are equivalent. the minimal DFA has states 0, 1/2, 3, and emptyset (dead state). The transition function of the minimal DFA is:

$\delta(0, a) = 1/2$
 $\delta(0, b) = \text{empty set (dead state)}$
 $\delta(1/2, a) = 3$
 $\delta(1/2, b) = 1/2$
 $\delta(3, a) = 1/2$
 $\delta(3, b) = 1/2$

III. Give a regular expression for the language of all strings over $\{a, b\}$ with three consecutive a 's.

Your answer:

$(a+b)^*aaa(a+b)^*$

IV. Let $L = \{a^n b^n \mid n \geq 0\}$. Design a PDA which accepts L .

Your answer:

Three states, 0, 1, and 2 2 is the only final state. The bottom of stack

symbol is z .

There is a self-loop at state 0. The labels on that loop are:

$z/a/az$, $a/a/aa$.

There is an arc from 0 to 1. The labels on that arc are:

$z/\lambda/z$, $a/\lambda/a$.

There is a self-loop at state 1. The label on that loop is:

$a/b/\lambda$.

There is an arc from 1 to 2. The label on that arc is

$z/\lambda/\lambda$.

V. Consider the following CNF grammar:

$S \rightarrow AB$

$A \rightarrow a$

$B \rightarrow b$

Use the CYK algorithm to prove that ab is in the language generated by the grammar.

Your answer:

$V[1,1] = A$, $V[2,2] = B$, $V[1,2] = S$

Therefore ab is generated by the grammar.

VI. Write exponents using the caret symbol. That is, a^2 should be written a^2

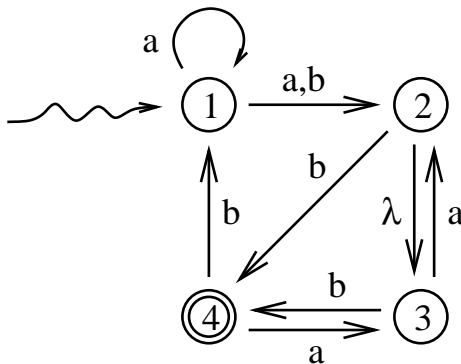
Practice Exam

1. True or False. T = true, F = false, and O = open, meaning that the answer is not known science at this time. In the questions below, \mathcal{P} and \mathcal{NP} denote \mathcal{P} -TIME and \mathcal{NP} -TIME, respectively.
 - (i) ____ Let L be the language over $\{a, b, c\}$ consisting of all strings which have more a 's than b 's and more b 's than c 's. There is some PDA that accepts L .
 - (ii) ____ The language $\{a^n b^n \mid n \geq 0\}$ is context-free.
 - (iii) ____ The language $\{a^n b^n c^n \mid n \geq 0\}$ is context-free.
 - (iv) ____ The language $\{a^i b^j c^k \mid j = i + k\}$ is context-free.
 - (v) ____ The intersection of any three regular languages is regular.
 - (vi) ____ The intersection of any two context-free languages is context-free.
 - (vii) ____ If L is a language and L^* is regular, then L must be regular.
 - (viii) ____ If L is a context-free language over an alphabet with just one symbol, then L is regular.
 - (ix) ____ The set of strings that your high school algebra teacher would accept as legitimate expressions is a context-free language.

- (x) ——— Every language accepted by a non-deterministic machine is accepted by some deterministic machine.
- (xi) ——— Every context-free language is generated by some unambiguous context-free grammar.
- (xii) ——— The problem of whether a given string is generated by a given context-free grammar is decidable.
- (xiii) ——— If G is a context-free grammar, the question of whether $L(G) = \emptyset$ is decidable.
- (xiv) ——— Every language generated by an unambiguous context-free grammar is accepted by some DPDA.
- (xv) ——— The language $\{a^n b^n c^n d^n \mid n \geq 0\}$ is decidable.
- (xvi) ——— Every problem that can be mathematically defined has an algorithmic solution.
- (xvii) ——— $\mathcal{P} = \mathcal{NP}$.
- (xviii) ——— There exists a polynomial time algorithm which finds the factors of any positive integer, where the input is given as a binary numeral.
- (xix) ——— The language consisting of all strings over $\{a, b\}$ which have more a 's than b 's is context-free.
- (xx) ——— Every context-free language is in \mathcal{P} .
- (xxi) ——— Every context-sensitive language is in \mathcal{P} .
- (xxii) ——— Every language generated by a general grammar is decidable.
- (xxiii) ——— The problem of whether two given context-free grammars generate the same language is decidable.
- (xxiv) ——— Every bounded function is recursive (that is, computable).
- (xxv) ——— Recall that if \mathcal{L} is a class of languages, $\text{co-}\mathcal{L}$ is defined to be the class of all languages that are not in \mathcal{L} . Let \mathcal{RE} be the class of all recursively enumerable languages. If L is in \mathcal{RE} and also L is in $\text{co-}\mathcal{RE}$, then L must be decidable.
- (xxvi) ——— If a language L is both \mathcal{NP} and $\text{co-}\mathcal{NP}$, L must be \mathcal{P} .
- (xxvii) ——— Every language is enumerable. That means, either L is finite or there is a one-to-one function from the positive integers to L .
- (xxviii) ——— There is a non-recursive function which grows faster than any recursive function. (Recursive function means computable function.)
- (xxix) ——— There exists a machine¹ that runs forever and outputs the string of decimal digits of π (the well-known ratio of the circumference of a circle to its diameter).

¹As always in automata theory, “machine” means abstract machine, a mathematical object whose memory and running time are **not** constrained by the size and lifetime of the known (or unknown) universe, or any other physical laws. A computer program, which can have any length, can be considered to be a machine; by assuming it runs on an abstract machine. If we want to discuss the kind of machine that exists (or could exist) physically, we call it a “physical machine.”

- (xxx) ——— For every real number x , there exists a machine that runs forever and outputs the string of decimal digits of x .
- (xxxii) ——— There is a polynomial time algorithm which determines whether any two regular expressions are equivalent.
- (xxxiii) ——— If two regular expressions are equivalent, there is a polynomial time proof that they are equivalent.
- (xxxiiii) ——— The halting problem is undecidable.
- Let L be the language consisting of all strings over the binary alphabet whose last three symbols are '010.' Design an NFA with four states which accepts L .
 - Let L be the language consisting of all strings over $\{a, b\}$ which do not contain the substring aab . Write a regular expression for L and draw a minimal DFA which accepts L . (Hint: 3 states.)
 - Consider the NFA shown below.
 - Draw a state diagram for a minimal DFA equivalent to that NFA.
 - Write a regular expression which describes the language accepted by that NFA.
 - Give a regular grammar which generates the language accepted by that NFA.



- Consider the language L generated by the CF grammar given below.

$$S \rightarrow wS$$

$$S \rightarrow iS$$

$$S \rightarrow iSeS$$

$$S \rightarrow a$$
 - Give a Chomsky Normal Form grammar for L .
 - Use the CYK algorithm to decide whether $iiwaea \in L$. Show all values of $V_{i,j}$.
- Let L be the set of all algebraic expressions, where the only operations permitted are addition, subtraction, and multiplication, where there are only two variables, x and y and no constants, and where multiplication is indicated by concatenation, as in the expression $x + y(xy + x)$. Parentheses can be used. Write a context-free grammar for L . Your grammar may be ambiguous.
- Design a PDA which accepts the Dyck language. To make it easier to grade, use a and b for left and right parentheses.