

## Computer Science 456/656 Spring 2021 Practice Examination April 8, 2021

The entire examination is 655 points. The actual exam will be shorter.

1. True or False. [5 points each] T = true, F = false, and O = open, meaning that the answer is not known to science at this time.
  - i ..... Let  $L$  be the language over  $\Sigma = \{a, b\}$  consisting of all strings of the form  $a^m b^n$ , where  $m, n \geq 0$ . Then  $L$  is a regular language.
  - ii ..... The complement of every regular language is regular.
  - iii ..... The Kleene closure of every context-free language is context-free.
  - iv ..... If a language has an ambiguous context-free grammar, then it is not accepted by any deterministic push-down automaton.
  - v ..... There is a PDA that accepts all valid C++ programs.
  - vi ..... The intersection of any two regular languages is regular.
  - vii ..... The intersection of any two context-free languages is context-free.
  - viii ..... The set of all base 7 numerals for positive integers  $n$  such that  $n \% 3 = 2$  is regular.
  - ix ..... Let  $L$  be the language over  $\Sigma = \{a, b\}$  consisting of all strings of the form  $a^m b^n c^m$ , where  $m, n \geq 0$ . Then  $L$  is a context-free language.
  - x ..... Let  $L$  be the language over  $\Sigma = \{a, b\}$  consisting of all strings of the form  $a^m b^n$ , where  $m \geq n$ . Then  $L$  is a context-free language.
  - xi ..... The complement of every context-free language is context-free.
  - xii ..... The union of any two context-free languages is context-free.
  - xiii ..... If a language has a context-free grammar, then it is accepted by some push-down automaton.
  - xiv ..... Every context-free language has an unambiguous context-free grammar.
  - xv ..... Every language that has an unambiguous context-free grammar is accepted by some DPDA.
  - xvi ..... Every deterministic machine is a non-deterministic machine.
  - xvii ..... The language consisting of all base 2 numerals for integer powers of 2 is regular.
  - xviii ..... There is a DPDA that accepts the language of all palindromes over the binary alphabet  $\{0, 1\}$ .
  - xix ..... The language  $\{a^n b^n c^n d^n \mid n \geq 0\}$  is recursive.
  - xx ..... The problem of whether a given context-free grammar generates all strings is decidable.

- xxi \_\_\_\_\_ The language  $\{a^i b^j c^k \mid j \geq i + k\}$  is context-free.
- xxii \_\_\_\_\_ If a language  $L$  is undecidable, there is no machine that enumerates  $L$  in canonical order.
- xxiii If  $\mathcal{L}$  is a class of languages,  $\text{co-}\mathcal{L}$  is defined to be the class of all languages that are not in  $\mathcal{L}$ .  
 \_\_\_\_\_ Let  $\mathcal{RE}$  be the class of all recursively enumerable languages. If  $L$  is in  $\mathcal{RE}$  and also  $L$  is in  $\text{co-}\mathcal{RE}$ , then  $L$  must be decidable.
- xxiv \_\_\_\_\_ The context-free grammar equivalence problem is in the class  $\mathcal{RE}$ .
- xxv \_\_\_\_\_ The context-free grammar equivalence problem is in the class  $\text{co-}\mathcal{RE}$ .
- xxvi \_\_\_\_\_ Every bounded function is recursive.
- xxvii \_\_\_\_\_ If  $P$  is a mathematical proposition that can be stated using  $n$  binary bits, and  $P$  has a proof, then  $P$  must have a proof whose length is  $O(2^{2^n})$ .
- xxviii \_\_\_\_\_ The complement of every recursive language is recursive.
- xxix \_\_\_\_\_ The complement of every recursively enumerable language is recursively enumerable.
- xxx \_\_\_\_\_ Every language which is generated by an general grammar is recursively enumerable.
- xxxi \_\_\_\_\_ The context-free grammar equivalence problem is undecidable.
- xxxii \_\_\_\_\_ Given any context-free grammar  $G$  and any string  $w \in L(G)$ , there is always a unique leftmost derivation of  $w$  using  $G$ .
- xxxiii \_\_\_\_\_ For any deterministic finite automaton, there is always a unique minimal non-deterministic finite automaton equivalent to it.
- xxxiv \_\_\_\_\_ Using multi-processors and other advanced technology, it is possible to design a machine which decides the halting problem.
- xxxv \_\_\_\_\_ The question of whether two regular expressions are equivalent is  $\mathcal{NP}$ -complete.
- xxxvi \_\_\_\_\_ The intersection of any context-free language with any regular language is context-free.
- xxxvii \_\_\_\_\_ The halting problem is recursively enumerable.
- xxxviii \_\_\_\_\_ The complement of every context-free language is context-free.
- xxxix \_\_\_\_\_ No language which has an ambiguous context-free grammar can be accepted by a DPDA.
- xl \_\_\_\_\_ The union of any two context-free languages is context-free.
- xli \_\_\_\_\_ The question of whether a given Turing Machine halts with empty input is decidable.
- xl ii \_\_\_\_\_ The class of languages accepted by non-deterministic Turing Machines is the same as the class of languages accepted by Turing Machines.
- xl iii \_\_\_\_\_ Let  $F(0) = 1$ , and let  $F(n) = 2^{F(n-1)}$  for  $n > 0$ . Then  $F$  is recursive.

- xliv ----- Every language which is accepted by some non-deterministic machine is accepted by some deterministic machine.
- xlv ----- The language of all regular expressions over the binary alphabet is a regular language.
- xlvi ----- There is no computer program that decides whether any two C++ programs are equivalent.
- xlvii ----- Every function that can be mathematically defined is recursive.
- xlviii ----- The language of all binary numerals for multiples of 23 is regular.
- xliv ----- The language of all binary strings which are the binary numerals for prime numbers is context-free.
- l ----- Every bounded function from integers to integers is Turing-computable. (We say that  $f$  is *bounded* if there is some  $B$  such that  $|f(n)| \leq B$  for all  $n$ .)
- li ----- The language of all palindromes over  $\{0, 1\}$  is inherently ambiguous.
- lii ----- The language  $\{a^i b^j c^k : i = j \text{ or } j = k\}$  is context-free, but is inherently ambiguous.
- liii ----- Every context-free grammar can be parsed by some deterministic top-down parser.
- liv ----- Every context-free grammar can be parsed by some non-deterministic top-down parser.
- lv ----- Commercially available parsers cannot use the LALR technique, since most modern programming languages are not context-free.
- lvi ----- The diagonal language is  $\mathcal{RE}$ .
- lvii ----- The diagonal language is  $\text{co-}\mathcal{RE}$ .
- lviii ----- The regular grammar membership problem is in  $\mathcal{NC}$ .
- lix ----- The context-free grammar membership problem is in  $\mathcal{NC}$ .
- lx ----- There is a polynomial time reduction of SAT to 3-CNF-SAT.
- lxi ----- There is a language which is neither  $\mathcal{RE}$  nor  $\text{co-}\mathcal{RE}$ .
- lxii -----  $\mathcal{P}\text{-SPACE} = \text{co-}\mathcal{P}\text{-SPACE}$
- lxiii ----- Regular expression equivalence is  $\mathcal{P}\text{-SPACE}$  complete.

2. Fill in the blanks.

- (a) If there is an easy reduction from  $L_1$  to  $L_2$ , then ----- is at least as hard as -----.
- (b) If a language is accepted by some Turing machine, it is ----- enumerable.
- (c) If  $L_1$  is  $\mathcal{NP}$  and  $L_2$  is  $\mathcal{NP}$ -complete, there must be a  $\mathcal{P}$ -TIME reduction of ----- to -----.
- (d) The language HALT is generated by an ----- grammar.

3. [30 points] State the Church-Turing thesis, and explain (in about 5 lines or less) why it is important.

4. [30 points] We give an LALR parser for the context-free grammar with start symbol  $S$  and the given productions. (Think of  $b$  and  $n$  as begin and end, or as  $\{$  and  $\}$ .) Walk through a computation of that parser for the input string  $bwwaiaebnn$ . The first configuration is  $\$_0 : bwwaiaebnn\$$ .

$$1 \ S \rightarrow a_2$$

$$2 \ S \rightarrow i_3 S_4$$

$$3 \ S \rightarrow i_3 S_4 e_5 S_6$$

$$4 \ S \rightarrow w_7 S_8$$

$$5 \ S \rightarrow b_9 L_{10} n_{11}$$

$$6 \ L \rightarrow L_{10} S_{12}$$

$$7 \ L \rightarrow \lambda$$

	a	i	e	w	b	n	\$	S	L
0	s2	s3		s7	s9			1	
1							<b>halt</b>		
2	r1	r1	r1	r1	r1	r1	r1		
3	s2	s3		s7	s9			4	
4	r2	r2	s5	r2	r2	r2	r2		
5	s2	s3		s7	s9			6	
6	r3	r3	r3	r3	r3	r3	r3		
7	s2	s3		s7	s9			8	
8	r4	r4	r4	r4	r4	r4	r4		
9	r7	r7		r7	r7	r7			10
10	s2	s3		s7	s9	s11		12	
11	r5	r5	r5	r5	r5	r5	r5		
12	r6	r6	r6	r6	r6	r6			



13. [20 points] Let  $\Sigma = \{0, 1\}$ , the binary alphabet. We say a string  $w$  over  $\Sigma$  is *mostly positive* if  $\#_1(w) > \#_0(w)$ . Let  $L$  be the set of mostly positive strings over  $\Sigma$ .

Give a context-free grammar for  $L$ . **Very hard.**

14. [10 points] What is a **reduction** of a language  $L_1$  to a language  $L_2$ ?

15. (a) [10 points] State the pumping lemma for regular languages accurately. If you have all the right words but in the wrong order, that means you truly do not understand the lemma, and you might get no partial credit at all.

- (b) [10 points] State the pumping lemma for context-free languages accurately. If you have all the right words but in the wrong order, that means you truly do not understand the lemma, and you might get no partial credit at all.
16. [20 points] Prove that a recursively enumerable language is accepted by some machine.
17. [20 points] Prove that a language is accepted by some machine if it is recursively enumerable.

18. [20 points] Prove that any language that can be recursively enumerated in canonical order is recursive.

19. [20 points] Prove that a recursive language can be recursively enumerated in canonical order.

20. [20 points] Give a polynomial time reduction of the subset sum problem to the partition problem.



21. [20 points] Give a polynomial time reduction of 3-CNF-SAT to the independent set problem. You may use figures in your explanation.