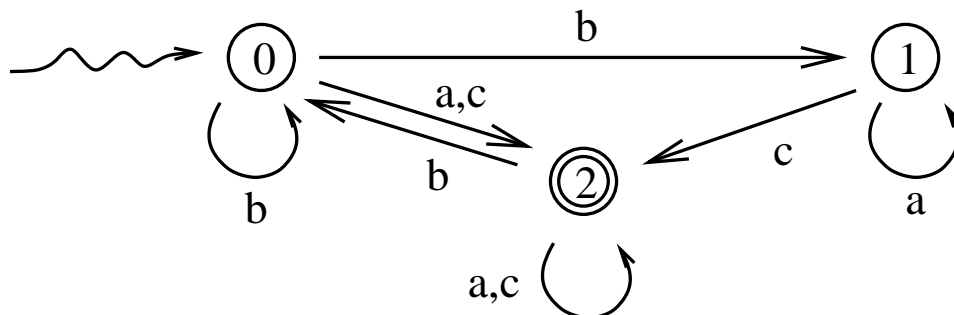


# University of Nevada, Las Vegas Computer Science 456/656 Spring 2021

## Practice Problems for the Examination on February 8, 2023

1. Review answers to homework1:  
<http://web.cs.unlv.edu/larmore/Courses/CSC456/S23/Assignments/hw1ans.pdf>
2. Review answers to homework2:  
<http://web.cs.unlv.edu/larmore/Courses/CSC456/S23/Assignments/hw2ans.pdf>
3. Find a minimal DFA equivalent to the NFA shown below.



4. True or False. If the question is currently open, write “O” or “Open.”
  - (i) ----- The complement of every regular language is regular.
  - (ii) ----- The complement of every context-free language is context-free.
  - (iii) ----- The complement of any  $\mathcal{P}$ -TIME language is  $\mathcal{P}$ -TIME.
  - (iv) ----- ~~The complement of any  $\mathcal{NP}$  language is  $\mathcal{NP}$ .~~
  - (v) ----- ~~The complement of any  $\mathcal{P}$ -SPACE language is  $\mathcal{P}$ -SPACE.~~
  - (vi) ----- The complement of every recursive language is recursive.
  - (vii) ----- The complement of every recursively enumerable language is recursively enumerable.
  - (viii) ----- ~~Every language which is generated by a general grammar is recursively enumerable.~~
  - (ix) ----- ~~The context-free membership problem is undecidable.~~
  - (x) ----- The factoring problem, where inputs are written in binary notation, is  $\text{co-}\mathcal{NP}$ .
  - (xi) ----- ~~If  $L_1$  reduces to  $L_2$  in polynomial time, and if  $L_2$  is  $\mathcal{NP}$ , and if  $L_1$  is  $\mathcal{NP}$ -complete, then  $L_2$  must be  $\mathcal{NP}$ -complete.~~
  - (xii) ----- Given any context-free grammar  $G$  and any string  $w \in L(G)$ , there is always a unique leftmost derivation of  $w$  using  $G$ .
  - (xiii) ----- For any deterministic finite automaton, there is always a unique minimal non-deterministic finite automaton equivalent to it.
  - ~~The question of whether two regular expressions are equivalent is known to be  $\mathcal{NP}$ -complete.~~

- (xiv) ~~----- The halting problem is recursively enumerable.~~
- (xv) ~~----- The union of any two context-free languages is context-free.~~
- (xvi) ~~----- The question of whether a given Turing Machine halts with empty input is decidable.~~
- (xvii) ~~----- The class of languages accepted by non-deterministic finite automata is the same as the class of languages accepted by deterministic finite automata.~~
- (xviii) ~~----- The class of languages accepted by non-deterministic push-down automata is the same as the class of languages accepted by deterministic push-down automata.~~
- (xix) ~~----- The intersection of any two context-free languages is context-free.~~
- (xx) ~~----- If  $L_1$  reduces to  $L_2$  in polynomial time, and if  $L_2$  is  $\mathcal{NP}$ , then  $L_1$  must be  $\mathcal{NP}$ .~~
- (xxi) ~~----- The language of all regular expressions over the binary alphabet is a regular language.~~
- (xxii) ~~----- Let  $\pi$  be the ratio of the circumference of a circle to its diameter. The problem of whether the  $n^{\text{th}}$  digit of the decimal expansion of  $\pi$  for a given  $n$  is equal to a given digit is decidable.~~
- (xxiii) ~~----- There cannot exist any computer program that can decide whether any two C++ programs are equivalent.~~
- (xxiv) ~~----- An undecidable language is necessarily  $\mathcal{NP}$ -complete.~~
- (xxv) ~~----- Every context-free language is in the class  $\mathcal{P}$ -TIME.~~
- (xxvi) ~~----- Every regular language is in the class  $\mathcal{NC}$ .~~
- (xxvii) ~~----- Every Function that can be mathematically defined is recursive.~~
- (xxviii) ~~----- The language of all binary strings which are the binary numerals for prime numbers is context-free.~~
- (xxix) ~~----- The language of all binary strings which are the binary numerals for prime numbers is regular.~~
- (xxx) ~~----- Every bounded function from integers to integers is Turing-computable. (We say that  $f$  is bounded if there is some  $B$  such that  $|f(n)| \leq B$  for all  $n$ .)~~
- (xxxi) ~~----- The language of all palindromes over  $\{0, 1\}$  is inherently ambiguous.~~
- (xxxii) ~~----- Every context-free grammar can be parsed by some deterministic top-down parser.~~
- (xxxiii) ~~----- Every context-free grammar can be parsed by some non-deterministic top-down parser.~~
- (xxxiv) ~~----- Commercially available parsers cannot use the LALR technique, since most modern programming languages are not context-free.~~
- (xxxv) ~~----- The boolean satisfiability problem is undecidable.~~
- (xxxvi) ~~----- If anyone ever proves that  $\mathcal{P} = \mathcal{NP}$ , then all one-way encoding systems will be insecure.~~
- (xxxvii) ~~----- If a string  $w$  is generated by a context-free grammar  $G$ , then  $w$  has a unique leftmost derivation if and only if it has a unique rightmost derivation.~~

- (xxxviii) ----- ~~A language  $L$  is in  $\mathcal{NP}$  if and only if there is a polynomial time reduction of  $L$  to SAT.~~
- (xxxix) ----- Every subset of a regular language is regular.
- (xl) ----- ~~The intersection of any context-free language with any regular language is context-free.~~
- (xli) ----- ~~Every language which is generated by a general grammar is recursively enumerable.~~
- (xlii) ----- ~~The question of whether two context-free grammars generate the same language is undecidable.~~
- (xliii) ----- ~~There exists some proposition which is true but which has no proof.~~
- (xliv) ----- The set of all binary numerals for prime numbers is in the class  $\mathcal{P}$ .
- (xlv) ----- ~~If  $L_1$  reduces to  $L_2$  in polynomial time, and if  $L_2$  is  $\mathcal{NP}$ , and if  $L_1$  is  $\mathcal{NP}$ -complete, then  $L_2$  must be  $\mathcal{NP}$ -complete.~~
- (xlvi) Given any context-free grammar  $G$  and any string  $w \in L(G)$ , there is always a unique leftmost derivation of  $w$  using  $G$ .
- (xlvii) ----- For any deterministic finite automaton, there is always a unique minimal non-deterministic finite automaton equivalent to it.
- (xlviii) ----- ~~The question of whether two regular expressions are equivalent is  $\mathcal{NP}$ -complete.~~
- (xlix) ----- ~~No language which has an ambiguous context-free grammar can be accepted by a DPDA.~~
  - (l) ----- The class of languages accepted by non-deterministic push-down automata is the same as the class of languages accepted by deterministic push-down automata.
  - (li) ----- The intersection of any two regular languages is regular.
  - (lii) ----- The intersection of any two context-free languages is context-free.
  - (liii) ----- If  $L_1$  reduces to  $L_2$  in polynomial time, and if  $L_2$  is  $\mathcal{NP}$ , then  $L_1$  must be  $\mathcal{NP}$ .
  - (liv) ----- Let  $F(0) = 1$ , and let  $F(n) = 2^{F(n-1)}$  for  $n > 0$ . Then  $F$  is recursive.
  - (lv) ----- Every language which is accepted by some non-deterministic machine is accepted by some deterministic machine.
  - (lvi) ----- The language of all regular expressions over the binary alphabet is a regular language.
  - (lvii) ----- Let  $\pi$  be the ratio of the circumference of a circle to its diameter. (That's the usual meaning of  $\pi$  you learned in school.) The problem of whether the  $n^{\text{th}}$  digit of  $\pi$ , for a given  $n$ , is equal to a given digit is decidable.
  - (lviii) ----- There cannot exist any computer program that decides whether any two given C++ programs are equivalent.
  - (lix) ----- An undecidable language is necessarily  $\mathcal{NP}$ -complete.
  - (lx) Every context-free language is in the class  $\mathcal{P}$ -TIME.

- (lxi) ----- Every function that can be mathematically defined is recursive.
- (lxii) ----- Every bounded function from integers to integers is Turing-computable. (We say that  $f$  is bounded if there is some  $B$  such that  $|f(n)| \leq B$  for all  $n$ .)
- (lxiii) Every context-free language is in the class  $\mathcal{P}$ -TIME.
- (lxiv) ----- Every function that can be mathematically defined is recursive.
- (lxv) ----- The language of all binary strings which are the binary numerals for multiples of 23 is regular.
- (lxvi) ----- ~~Every bounded function from integers to integers is Turing-computable. (We say that  $f$  is bounded if there is some  $B$  such that  $|f(n)| \leq B$  for all  $n$ .)~~
- (lxvii) Commercially available parsers cannot use the LALR technique, since most modern programming languages are not context-free.
- (lxviii) ----- ~~If anyone ever proves that  $\mathcal{P} = \mathcal{NP}$ , then all public key/private key encryption systems will be known to be insecure.~~
- (lxix) ----- ~~If a sequence of fractions converges to a real number  $x$ , then  $x$  must be a recursive real number.~~  
----- ~~If a machine outputs a sequence of fractions which converges to a real number  $x$ , then  $x$  must be a recursive real number.~~