## University of Nevada, Las Vegas Computer Science 456/656 Spring 2024 <br> CSC 456/656 Fall 2024 Answers to Third Examination April 10, 2024

Name:
The entire test is 365 points.
A binary function is defined to be a function $F$ on binary strings such that, for each binary string $w, F(w)$ is a binary string. (Of course, the strings could be numerals.)

1. True or False.[5 points each] $\mathrm{T}=$ true, $\mathrm{F}=$ false, and $\mathrm{O}=$ open, meaning that the answer is not known science at this time.
(i) $\mathbf{T}$ Every language is a subset of some regular language.
(ii) O If $L$ is any $\mathcal{P}$-TIME language, there is an $\mathcal{N C}$ reduction of the Boolean circuit problem (CVP) to $L$.
(iii) $\mathbf{T}$ The concatenation of any two RE languages is RE.
(iv) $\mathbf{F}$ Every function that can be mathematically defined is bounded by some recursive function.
(v) $\mathbf{T}$ There are uncountably many languages over the binary alphabet.
(vi) $\mathbf{F}$ There are uncountably many RE languages over the binary alphabet.
(vii) $\mathbf{T}$ There is an algorithm which determines whether a given list of $n$ positive integers has a sublist whose total is a given number $S$.
(viii) $\mathbf{T}$ If $L_{1}$ is $\mathcal{N} \mathcal{P}$-complete and $L_{2}$ is $\mathcal{N P}$ and there is a $\mathcal{P}$-TIME reduction of $L_{1}$ to $L_{2}$, then $L_{2}$ must be $\mathcal{N} \mathcal{P}$-complete.
(ix) T Nick's Class is closed under Kleene closure.
(x) T Any computation that can be done by any machine can be done by some Pascal program.
(xi) $\mathbf{T}$ Multiplication of integer matrices is $\mathcal{N C}$.
(xii) $\mathbf{T}$ There is an $\mathcal{P}$-SPACE algorithm which decides SAT.
(xiii) O Every dynamic program problem can be worked by polynomially many processors in polylogarithmic time.
(xiv) $\mathbf{F}$ Let $L$ be any undecidable $\mathcal{R E}$ language, and let $M_{L}$ be a machine which accepts $L$. For any string $w \in M$, let $F_{L}(w)$ be the number of steps $M_{L}$ executes, if its input is $w$. Now define $T_{L}(n)=\left\{\max F_{L}(w): w \in L\right.$ and $\left.|w|=n\right\}$. Then $T_{L}$ is recursive.
(xv) $\mathbf{O}$ There is a polynomial time reduction of the subset sum problem to 2-SAT.
(xvi) $\mathbf{T}$ Every $\mathcal{P}$-Time problem has an $\mathcal{N C}$ reduction to the Circuit Value Problem.
(xvii) $\mathbf{T}$ If a language $L$ is accepted by a non-deterministic machine, then $L$ must be accepted by some deterministice machine.
(xviii) $\mathbf{T}$ Every $\mathcal{N C}$ language is context-free.
(xix) $\mathbf{T}$ The language $\left\{a^{n} b^{n} c^{n} d^{n}: n \geq 0\right\}$ is $\mathcal{N C}$.
(xx) F The set of all languages over the binary alphabet is countable.
(xxi) $\mathbf{T}$ The context-free grammar equivalence problem is co- $\mathcal{R E}$.
(xxii) $\mathbf{T}$ The set of all binary numerals for prime numbers is $\mathcal{P}$-TIME.
(xxiii) $\mathbf{T}$ If $L$ is a context-free language over the unary alphabet, then $L$ must be regular.
(xxiv) $\mathbf{F}$ The union of any two undecidable languages is undecidable.
(xxv) $\mathbf{T}$ co- $\mathcal{P}$-TIME $=\mathcal{P}$-TIME.
(xxvi) $\mathbf{O}$ There exists a one-way function.
(xxvii) $\mathbf{T}$ The complement of any $\mathcal{P}$-SPACE language is $\mathcal{P}$-SPACE.
(xxviii) $\mathbf{T}$ The jigsaw puzzle problem is $\mathcal{N P}$ complete.
(xxix) $\mathbf{T}$ The furniture mover's problem is $\mathcal{P}$-SPACE complete.
(xxx) T The complement of any undecidable language is undecidable.
(xxxi) T If a Boolean expression is satisfiable, there is a polynomial time proof that it is satisfiable.
(xxxii) $\mathbf{T}$ If there is a recursive reduction from the halting problem to $L$, then $L$ must be undecidable.
(xxxiii) $\mathbf{F}$ If $L$ is undecidable, there must be a recursive reduction from the halting problem to $L$.
2. Fill in the blanks.
(a) [5 points] If $L_{1}$ is $\mathcal{N} \mathcal{P}$-complete and $L_{2}$ is $\mathcal{N P}$, and there is a polynomial time reduction of $L_{1}$ to $L_{2}$, then $L_{2}$ must be $\mathcal{N} \mathcal{P}$-complete.
(b) [5 points] A language is decidable or recursive. if and only if it is both RE and co-RE.
3. Here is a list of problems or languages. For each problem, enter $\mathbf{T}$ if it is known to be $\mathcal{N} \mathcal{P}$-complete, $\mathbf{F}$ if it is not known to be $\mathcal{N} \mathcal{P}$-complete. [5 points each]
(a) $\mathbf{T}$ SAT
(b) F 2-SAT
(c) $\mathbf{T}$ 3-SAT
(d) $\mathbf{F}$ Rush Hour
(e) $\mathbf{F}$ The Boolean circuit problem.
(f) $\mathbf{F}$ Integer factoring, using binary numerals.
(g) $\mathbf{T}$ The tiling problem.
(h) $\mathbf{F}$ The furniture mover's problem.
(i) $\mathbf{T}$ The bin packing problem. Given a set of bins, each with a given capacity, and given a set of items, can all the items be packed into the bins?
4. [20 points] Explain how to find the maximum of a list of $n$ integers in logarithmic time using $n$ processors.
partition the list into pairs. If there is one left over, call it a pair. Find the maximum of each pair in $O(1)$ time, using $n$ processors in parallel, All these maxima form a new list half as long as the original. Repeat the process until there is only one number left. That will be the maximum of the original list.
5. [20 points] Give a definition of a recursive real number. (There is more than one correct definition.)

Each of these is a correct answer. A real number $s$ is recursive if:

- There is a program that runs forever, writing the decimal expansion of $x$.
- There is a program that, given an integer $n$, writes the $n^{\text {th }}$ digit of the decimal expansion of $x$.
- There is a program that, given a rational number $q$, decides whether $q<x$.
- There is a program that, given a positive integer $q$, writes the integer $p$ such that $\frac{p}{q} \leq x<\frac{p+1}{q}$

Those are not the only ones.
7. [20 points] State the pumping lemma for context-free languages.
8. [20 points] Let $G$ be the following context-free gammar with start symbol $E$. Stack states are indicated. An LALR parser for $G$ is given below. Give a complete computation of the parser if the input string is $x-x *-(-x+x)$.

1. $E \rightarrow E_{1,11}+{ }_{2} E_{3}$
2. $E \rightarrow E_{1,11}{ }_{4} E_{5}$
3. $E \rightarrow E_{1,3,5,11} *_{6} E_{7}$
4. $E \rightarrow-{ }_{8} E_{9}$
5. $E \rightarrow\left({ }_{10} E_{11}\right)_{12}$
6. $E \rightarrow x_{13}$

|  | $x$ | + | - | $*$ | $($ | $)$ | $\$$ | $E$ |
| ---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | s13 |  | s8 |  | s10 |  |  | 1 |
| 1 |  | s2 | s4 | s6 |  |  | halt |  |
| 2 | s13 |  | s8 |  | s10 |  |  | 3 |
| 3 |  | r1 | r1 | s6 |  | r1 | r1 |  |
| 4 | s13 |  | s8 |  | s10 |  |  | 5 |
| 5 |  | r2 | r2 | s6 |  | r2 | r2 |  |
| 6 | s13 |  | s8 |  | s10 |  |  | 7 |
| 7 |  | r3 | r3 | r3 |  | r3 | r3 |  |
| 8 | s13 |  | s8 |  | s10 |  |  | 9 |
| 9 |  | r4 | r4 | r4 |  | r4 | r4 |  |
| 10 | s13 |  | s8 |  | s10 |  |  | 11 |
| 11 |  | s2 | s4 | s6 |  | s12 |  |  |
| 12 |  | r5 | r5 | r5 |  | r5 | r5 |  |
| 13 |  | r6 | r6 | r6 |  | r6 | r6 |  |


| $\$_{0}$ | $x-x *-(-x+x)$ |  |  |
| :---: | :---: | :---: | :---: |
| $\$_{0} x_{13}$ | $-x *-(-x+x)$ |  | $s 13$ |
| $\$_{0} E_{1}$ | $-x *-(-x+x)$ | 6 | $r 6$ |
| $\$_{0} E_{1}{ }_{4}$ | $x *-(-x+x)$ | 6 | $s 4$ |
| $\$_{0} E_{1}-{ }_{4} x_{13}$ | * $-(-x+x)$ | 6 | $s 13$ |
| $\$_{0} E_{1}-{ }_{4} E_{5}$ | * $-(-x+x)$ | 66 | $r 6$ |
| $\$_{0} E_{1}-{ }_{4} E_{5} *_{6}-_{8}$ | $(-x+x)$ | 66 | $s 8$ |
| $\$_{0} E_{1}-{ }_{4} E_{5} *_{6}-{ }_{8}(10$ | $-x+x)$ | 66 | s10 |
| $\$_{0} E_{1}-{ }_{4} E_{5} *_{6}-{ }_{8}\left({ }_{10}-8\right.$ | $x+x)$ | 66 | $s 8$ |
| $\$_{0} E_{1}-{ }_{4} E_{5} *{ }_{6}-{ }_{8}\left({ }_{10}-{ }_{8} x_{13}\right.$ | +x) | 66 | s13 |
| $\$_{0} E_{1}-{ }_{4} E_{5} * 6-{ }_{8}\left({ }_{10}-{ }_{8} E_{9}\right.$ | +x) | 666 | $r 6$ |
| $\$_{0} E_{1}-{ }_{4} E_{5} * 6-{ }_{8}\left({ }_{10} E_{11}\right.$ | +x) | 6664 | $r 4$ |
| $\$_{0} E_{1}-_{4} E_{5} *_{6}-{ }_{8}\left({ }_{10} E_{11}+2\right.$ | $x)$ | 6664 | $s 2$ |
| $\$_{0} E_{1}-{ }_{4} E_{5} * 6-{ }_{8}\left({ }_{10} E_{11}+{ }_{2} x_{13}\right.$ | ) | 6664 | s12 |
| $\$_{0} E_{1}-{ }_{4} E_{5} * 6-{ }_{8}\left({ }_{10} E_{11}+{ }_{2} E_{3}\right.$ | ) | 66646 | $r 6$ |
| $\$_{0} E_{1}-{ }_{4} E_{5} * 6-{ }_{8}\left({ }_{10} E_{11}\right.$ | ) | 666461 | $r 1$ |
| $\$_{0} E_{1}-{ }_{4} E_{5} *_{6}-{ }_{8}\left({ }_{10} E_{11}\right)_{12}$ |  | 666461 | $s 12$ |
| $\$_{0} E_{1}-{ }_{4} E_{5} *_{6}-{ }_{8} E_{9}$ |  | 6664615 | $r 5$ |
| $\$_{0} E_{1}{ }_{4} E_{5} *_{6} E_{7}$ |  | 66646154 | $r 4$ |
| $\$_{0} E_{1}-{ }_{4} E_{5}$ |  | 666461543 | $r 3$ |
| $\$_{0} E_{1}$ |  | 6664615432 | $r 2$ |
| halt |  |  |  |

9. Consider the following well-known complexity classes:
$\mathcal{N C} \subseteq \mathcal{P}-$ TIME $\subseteq \mathcal{N} \mathcal{P} \subseteq \mathcal{P}$-SPACE $\subseteq \mathbf{E X P}$-TIME $\subseteq \mathbf{E X P}$-SPACE
(a) [5 points] Which of the above complexity classes is the smallest class which is known to contain SAT, the Boolean satisfiability problem? $\mathcal{N} \mathcal{P}$
(b) [5 points] Which of the above complexity classes is the smallest class which is known to contain the connectivity problem for graphs? $\mathcal{P}$-TIME
(c) [5 points] Which of the above complexity classes is the smallest class which is known to contain the context-free language membership problem? $\mathcal{N C}$
(d) [5 points] Which of the above complexity classes is the smallest class which is known to contain every sliding block problem? $\mathcal{P}$-SPACE
(e) [5 points] Which of the above complexity classes is the smallest class which is known to contain integer matrix multiplication? $\mathcal{N C}$
10. [20 points] Give the verifier definition of the class $\mathcal{N} \mathcal{P}$. A language $L$ is $\mathcal{N} \mathcal{P}$ if and only if there is some machine $M$ and some integer $k$ such that:
(a) For any string $w \in L$ there is a string $c$ such that $M$ accepts $(w, c)$ within $n^{k}$ steps, where $n=|w|$
(b) For any string $w \notin L M$ does not accept $(w, c)$ for any string $c$.
11. [20 points] Prove that the halting problem is undecidable.

The proof is in Tests/stdy3ans.pdf.

