

## $\mathcal{NC}$ Addition

We assume that we have two length  $n$  binary numerals,  $\langle x \rangle$  and  $\langle y \rangle$ . We can compute the binary numeral  $\langle x + y \rangle$  in  $O(\log n)$  time using  $O(n)$  processors; the computation is thus in the class  $\mathcal{NC}$ .

In our example,  $n = 32$ . The numerals  $\langle x \rangle$  and  $\langle y \rangle$  are shown in the first two rows of the matrix below; the leftmost bit is the 31<sup>st</sup> bit, while the rightmost bit is the 0<sup>th</sup> bit. The third row shows the **save** bits, the sequence of bits  $s_{31}, \dots, s_0$  obtained by adding bits of  $x$  and  $y$  modulo 2.

Our goal is to compute the 33 bit numeral  $\langle x + y \rangle$ . Let  $c_i$  be the carry bit from the  $(i - 1)$ <sup>st</sup> place to the  $i$ <sup>th</sup> place. Since there is no carry bit to the 0<sup>th</sup> place, we have  $c_0 = 0$ . The  $i$ <sup>th</sup> bit of  $\langle x + y \rangle$  is  $s_i + c_i \% 2$ .

The most difficult part of this problem is computing the carry bits  $c_{32}, c_{31}, \dots, c_1$  in logarithmic time. You might guess that that is impossible, since a change in the 0<sup>th</sup> place could effect all carry bits. Our trick is to consider both possibilities at each place simultaneously.

### Places and Blocks

We define  $\text{Place}[i]$  to be consist of the  $i$ <sup>th</sup> bits of  $x$  and  $y$ . In our example,  $\text{Place}[0] = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$ ,  $\text{Place}[1] = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$ , etc. For  $i \geq j$ , let  $\text{Block}[i, j] = \text{Place}[i]\text{Place}[i - 1] \cdots \text{Place}[j]$ .

Each block defines a function  $F[i, j]$  from  $\{0, 1\}$  to  $\{0, 1\}$ .  $\text{Block}[i, j]$  defines a function from  $c_j$  to  $c_{i+1}$ .  $\text{TypeBlock}[i, j] \in \{A, B, C\}$  as follows

$$\text{TypeBlock}[i, j] = \begin{cases} A & \text{if } c_{i+1} = 0 \\ B & \text{if } c_{i+1} = c_j \\ C & \text{if } c_{i+1} = 1 \end{cases}$$

We first write the type of each place in the third row of the table using the rules:

$$\begin{aligned} \text{Type}\begin{bmatrix} 0 \\ 0 \end{bmatrix} &= A \\ \text{Type}\begin{bmatrix} 0 \\ 1 \end{bmatrix} &= \text{Type}\begin{bmatrix} 1 \\ 0 \end{bmatrix} = B \\ \text{Type}\begin{bmatrix} 1 \\ 1 \end{bmatrix} &= C \end{aligned}$$

For example,  $\text{TypePlace}[5] = A$ ,  $\text{TypePlace}[4] = C$ , and  $\text{TypePlace}[3] = \text{TypePlace}[2] = B$ , We fill in the rest of the table using *type algebra*, defined by the matrix:

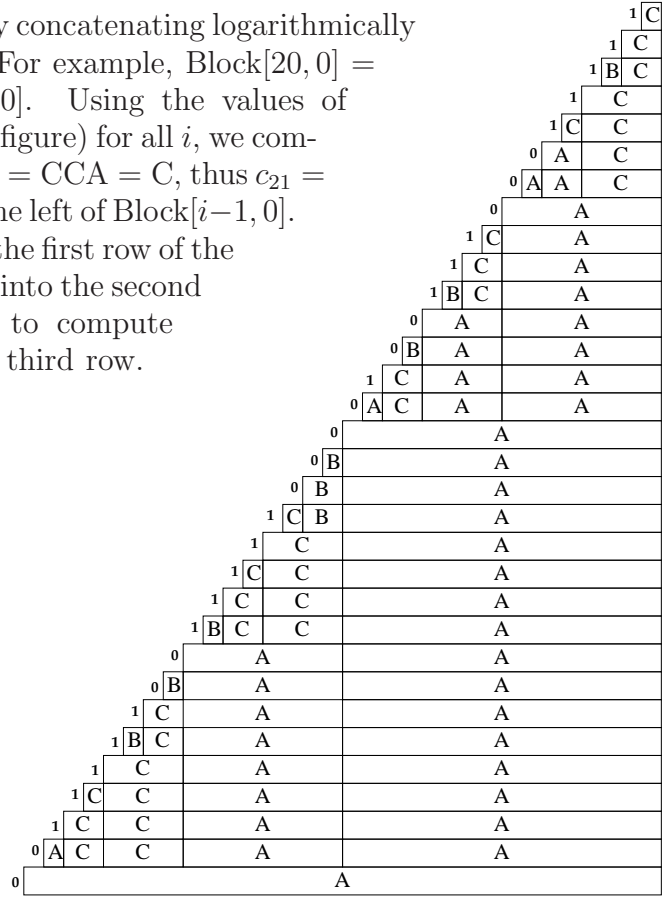
	A	B	C
A	A	A	A
B	A	B	C
C	C	C	C

In five steps, we compute the types of concatenations of blocks of length powers of 2. We compute the types of 16 blocks of size 2, then 8 blocks of size 4, then 4 blocks of size 8, then 2 blocks of size 16, and finally one block of size 32.

x	0	0	1	1	1	0	1	0	0	0	1	1	1	1	1	0	1	0	0	0	0	0	1	1	1	0	0	1	1	0	0	1	0	0	1
y	1	0	1	1	0	1	1	1	0	1	0	1	0	1	0	1	0	0	1	1	0	1	0	1	0	0	0	1	0	1	1	1			
save	1	0	0	0	1	1	0	1	0	1	1	0	1	0	1	1	1	0	1	1	0	1	1	0	1	0	0	0	1	1	1	0			
	B	A	C	C	B	B	C	B	A	B	B	C	B	C	B	B	B	A	B	B	A	B	B	C	B	A	A	C	B	B	B	C			
	A	C	B	C	A	C	C	B	A	B	A	C	A	A	B	C																			
	A		C		A		C		A		A		A		C																				
	A				A				A				A																						
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Since  $c_0 = 0$ , we have  $c_i = \begin{cases} 0 & \text{if TypeBlock}[i + 1, 0] \in \{A, B\} \\ 1 & \text{if TypeBlock}[i + 1, 0] = C \end{cases}$

We compute all  $\text{TypeBlock}[i-1, 0]$ , by concatenating logarithmically many blocks of size a power of 2. For example,  $\text{Block}[20, 0] = \text{Block}[20, 20] \text{Block}[19, 16] \text{Block}[15, 0]$ . Using the values of  $\text{TypeBlock}[i-1, 0]$  (not shown in the figure) for all  $i$ , we compute all  $c_i$ . For example,  $\text{Type}[20, 0] = \text{CCA} = C$ , thus  $c_{21} = 1$ . Each  $c_i$  is shown in the figure to the left of  $\text{Block}[i-1, 0]$ . Finally, we write the carry bits into the first row of the table below, then copy the save bits into the second row, and use addition modulo 2 to compute the bits of  $x + y$ , which are in the third row.



carries	0	0	1	1	1	1	1	0	0	1	1	1	1	1	0	0	0	0	1	0	0	1	1	1	0	0	0	1	1	1	1	1
save	1	0	0	0	1	1	0	1	0	1	1	0	1	0	1	1	1	0	1	1	0	1	1	0	1	0	0	0	1	1	1	0
x+y	0	1	1	1	1	0	0	0	1	1	0	0	1	0	0	1	1	1	1	1	0	0	0	1	0	1	1	0	0	0	0	