Exercises

0.1. In each of the following situations, indicate whether \( f = O(g) \), or \( f = \Omega(g) \), or both (in which case \( f = \Theta(g) \)).

<table>
<thead>
<tr>
<th>( f(n) )</th>
<th>( g(n) )</th>
</tr>
</thead>
<tbody>
<tr>
<td>( n - 100 )</td>
<td>( n - 200 )</td>
</tr>
<tr>
<td>( n^{1/2} )</td>
<td>( n^{2/3} )</td>
</tr>
<tr>
<td>( 100n + \log n )</td>
<td>( n + (\log n)^2 )</td>
</tr>
<tr>
<td>( n \log n )</td>
<td>( 10n \log 10n )</td>
</tr>
<tr>
<td>( \log 2n )</td>
<td>( \log 3n )</td>
</tr>
<tr>
<td>( 10 \log n )</td>
<td>( \log(n^2) )</td>
</tr>
<tr>
<td>( n^{1.01} )</td>
<td>( n \log^2 n )</td>
</tr>
<tr>
<td>( n^2 / \log n )</td>
<td>( n(\log n)^2 )</td>
</tr>
<tr>
<td>( n^{0.1} )</td>
<td>( (\log n)^{10} )</td>
</tr>
<tr>
<td>( (\log n)^{\log n} )</td>
<td>( n / \log n )</td>
</tr>
<tr>
<td>( \sqrt{n} )</td>
<td>( (\log n)^3 )</td>
</tr>
<tr>
<td>( n^{1/2} )</td>
<td>( 5^{\log_2 n} )</td>
</tr>
<tr>
<td>( n^{2n} )</td>
<td>( 3^n )</td>
</tr>
<tr>
<td>( 2^n )</td>
<td>( 2^{n+1} )</td>
</tr>
<tr>
<td>( n! )</td>
<td>( 2^n )</td>
</tr>
<tr>
<td>( (\log n)^{\log n} )</td>
<td>( 2(\log_2 n)^2 )</td>
</tr>
<tr>
<td>( \sum_{i=1}^{n} i^k )</td>
<td>( n^{k+1} )</td>
</tr>
</tbody>
</table>

0.2. Show that, if \( c \) is a positive real number, then \( g(n) = 1 + c + c^2 + \cdots + c^n \) is:

(a) \( \Theta(1) \) if \( c < 1 \).

(b) \( \Theta(n) \) if \( c = 1 \).

(c) \( \Theta(c^n) \) if \( c > 1 \).

The moral: in big-\( \Theta \) terms, the sum of a geometric series is simply the first term if the series is strictly decreasing, the last term if the series is strictly increasing, or the number of terms if the series is unchanging.

0.3. The Fibonacci numbers \( F_0, F_1, F_2, \ldots \), are defined by the rule

\[ F_0 = 0, \quad F_1 = 1, \quad F_n = F_{n-1} + F_{n-2}. \]

In this problem we will confirm that this sequence grows exponentially fast and obtain some bounds on its growth.

(a) Use induction to prove that \( F_n \geq 2^{0.5n} \) for \( n \geq 6 \).

(b) Find a constant \( c < 1 \) such that \( F_n \leq 2^{cn} \) for all \( n \geq 0 \). Show that your answer is correct.

(c) What is the largest \( c \) you can find for which \( F_n = \Omega(2^n) \)?

0.4. Is there a faster way to compute the \( n \)th Fibonacci number than by \texttt{fib2} (page 13)? One idea involves matrices.

We start by writing the equations \( F_1 = F_1 \) and \( F_2 = F_0 + F_1 \) in matrix notation:

\[
\begin{pmatrix}
F_1 \\
F_2
\end{pmatrix} =
\begin{pmatrix}
1 & 1 \\
0 & 1
\end{pmatrix}
\begin{pmatrix}
F_0 \\
F_1
\end{pmatrix}.
\]

Similarly,

\[
\begin{pmatrix}
F_3 \\
F_2
\end{pmatrix} =
\begin{pmatrix}
1 & 1 \\
0 & 1
\end{pmatrix}
\begin{pmatrix}
F_2 \\
F_3
\end{pmatrix} =
\begin{pmatrix}
1 & 1 \\
0 & 1
\end{pmatrix}^2
\begin{pmatrix}
F_0 \\
F_1
\end{pmatrix}.
\]