

University of Nevada, Las Vegas Computer Science 477/677 Fall 2021

Assignment 6: Due Monday November 1, 2021 11:59 pm

Name: _____

You are permitted to work in groups, get help from others, read books, and use the internet. Turn in the completed assignment on canvas, using instructions given to you by the grader, Mr. Heerdt, by 11:59 PM November 1.

“Pseudocode” means “false code.” It is supposed to represent a program in such a way to be humanly understandable, but need not follow any particular rules. Here is pseudocode for selection sort.

```
// Selection sort on an array A of size n.  Indices run from 1 to n.
For all i from 1 to n-1:  // select the ith smallest item
                        // place it in A[i]
For all j from i+1 to n: // replace A[i] by A[j] if it is smaller
  If(A[j] < A[i])
    Swap(A[i],A[j])      // we assume the code for Swap is available
```

In Problems we assume that G is a weighted directed graph with vertices numbered 0 through $n-1$. We let $E[i, j]$ be the weight of the arc from vertex i to vertex j . If no such arc exists, $E[i, j] = \infty$ by default.

1. Let G be defined by the following array of out-neighbor lists. If list i contains the entry (j,w) , there is an arc from i to j of weight w .

0 :	(1, 5)(3, -1)
1 :	(0, 1)(2, 4)
2 :	(0, 2)(3, 5)
3 :	(5, 3)(4, 1)
4 :	(5, 1)
5 :	(4, 2)

Fill in the array of in-neighbor lists of G .

0 :	
1 :	
2 :	
3 :	
4 :	
5 :	

Identify the strong components of G .

2. Write pseudocode for the Bellman-Ford algorithm to solve the single source minpath problem for G , where vertex 0 is the source. Your output should be two arrays of size n :
 1. $\text{distance}[i]$, the shortest path distance from the source to vertex i . If G has negative edges, the distance array could have negative entries. If no path from 0 to i exists, $\text{distance}[i]$ is infinity.
 2. $\text{back}[i]$, the next-to-the last vertex on the path of shortest length from the source to vertex i . Of course, $\text{back}[0]$ is undefined, and $\text{back}[i]$ is undefined if there is no path from the source to vertex i .

No minpath algorithm works if the directed graph has a “negative cycle,” that is, a cycle of negative weight. But, if the code is written properly, the Bellman-Ford algorithm can halt with an error message if G has a negative cycle. Try to do this in your code. If there is no negative cycle, the time complexity of your code should be $O(m\ell)$, where m is the number of arcs of G and ℓ is the maximum number of edges in any shortest length path.

3. Write pseudocode for the Floyd-Warshall algorithm to solve the all-pairs minpath problem for G . Assume there G has no negative cycle. The output of the algorithm is two square matrices:
 1. $\text{distance}[i][j]$, the least weight of any path from i to j .
 2. $\text{forward}[i][j]$, the second vertex of the least weight path from i to j . Note that $\text{forward}[i][i]$ is undefined.

4. Walk through Dijkstra's algorithm for the single source minpath problem for the directed graph illustrated below. Instead of numbering the vertices 0 through 19, I have assigned them letters from A to T. The source vertex is S.

After each iteration of the main loop, show

1. The array `dist`, where `dist[x]` is the smallest length of any path found so far from S to x. (Initially, `dist[x] = ∞` for most x.)
2. The array `back`, where `back[x]` is the next-to-the last vertex on the path of smallest weight found so far from S to x.
- 3 The contents of heap. Do not try to show the structure of the heap, simply list its members.

Attach an extra sheet of paper if necessary.

