

# Levenshtein Edit Distance

Given two strings  $u$  and  $v$  over an alphabet  $\Sigma$ , the Levenshtein distance from  $u$  to  $v$  (or from  $v$  to  $u$ ) is the number of edit steps needed to change  $u$  to  $v$ , where an edit step is one of the following:

1. Delete a symbol.
2. Insert a symbol.
3. Replace a symbol.

For example, the Levenshtein distance from WARM to BEAR is 3, since we can change WARM to BEAR with three edit steps:

WARM

1. delete M: WAR
2. replace W with B: BAR
3. insert E: BEAR

Levenshtein distance is computed using dynamic programming. Let  $n$  be the length of  $u$  and  $m$  the length of  $v$ . Let  $u[i]$  be the prefix of  $u$  of length  $i$  and let  $v[j]$  be the prefix of  $v$  of length  $j$ , for  $0 \leq i \leq n$  and  $0 \leq j \leq m$ .

Subproblem( $i,j$ ) is defined to be the computation of the Levenshtein distance from  $u[i]$  to  $v[j]$ , which we call  $L[i,j]$ . There are  $(n+1)(m+1)$  subproblems. The final answer is  $L[n,m]$ .

The program is as follows. Let  $u_i, v_i$  be the  $i^{\text{th}}$  symbol of  $u$  and the  $j^{\text{th}}$  symbol of  $v$ , respectively. The program is as follows:

For all  $i$  let  $L[i, 0] = i$

For all  $j$  let  $L[0, j] = j$

For all  $1 \leq i \leq n$

For all  $1 \leq j \leq m$

For all  $1 \leq j \leq m$

If( $u[i] = v[j]$ )

$L[i,j] = \min\{L[i-1,j]+1, L[i,j-1]+1, L[i-1,j-1]\}$

else

$L[i,j] = \min\{L[i-1,j]+1, L[i,j-1]+1, L[i-1,j-1]+1\}$

## Example

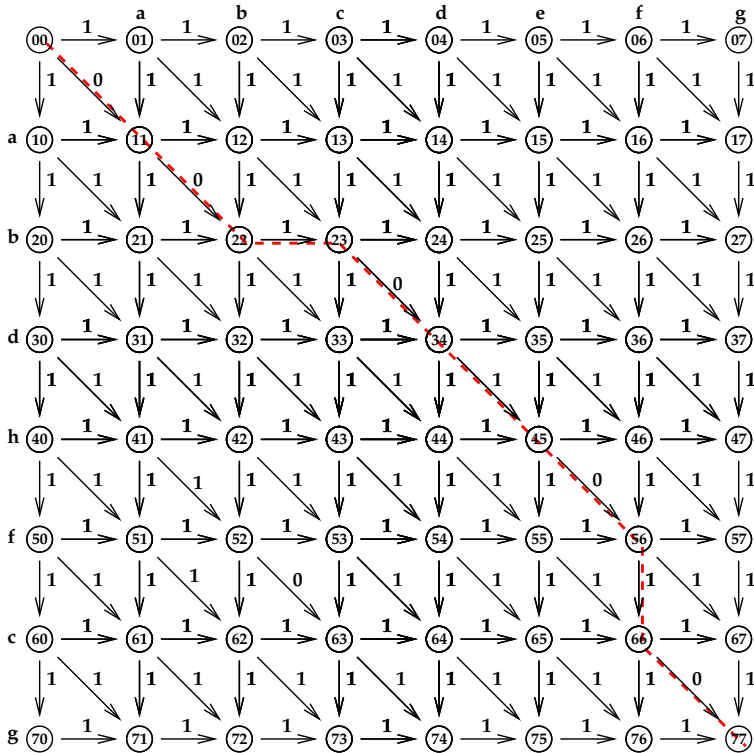
The following matrix shows the values of  $L$ .

			B	E	A	R
		0	1	2	3	4
	0	0	1	2	3	4
W	1	1	1	2	3	4
A	2	2	2	2	2	3
R	3	3	3	3	3	2
M	4	4	4	4	4	3

The Levenshtein distance is  $L[4, 4] = 3$

### Another Example

Compute the Levenshtein edit distance from `abdhhfbg` to `abcdefg`. The Levenshtein edit problem reduces to the single source shortest path problem. Finding the edit distance between strings of length  $m$  and  $n$  reduces to finding the shortest distance between two vertices in a weighted directed graph with  $(m + 1)(n + 1)$  vertices. We show that graph below, as well as the shortest path. The source is labeled 00, and the shortest path from 00 to  $ij$  is the value of  $L[i, j]$  in the pseudocode, which is the entry in the  $i^{\text{th}}$  row,  $j^{\text{th}}$  column of the matrix.



		a	b	c	d	e	f	g
	0	1	2	3	4	5	6	7
a	1	0	1	2	3	4	5	6
b	2	1	0	1	2	3	4	5
d	3	2	1	1	1	2	3	4
h	4	3	2	2	2	2	3	4
f	5	4	3	3	3	3	2	3
c	6	5	4	3	4	4	3	3
g	7	6	5	5	5	5	4	3

The Levenshtein distance is  $L[7, 7] = 3$ . The steps are:

1. insert c between b and d,
2. change h to e,
3. delete the c after f.