

University of Nevada, Las Vegas Computer Science 456/656 Fall 2025

Assignment 6: Due November 8, 2025, 11:59:59 PM

Name: _____

You are permitted to work in groups, get help from others, read books, and use the internet. You will receive a message from the graduate assistant, Sabrina Wallace, telling you how to turn in the assignment.

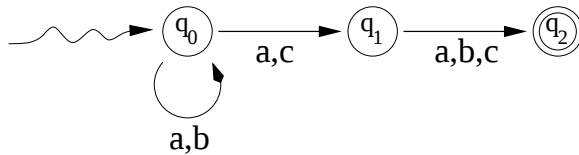
There are problems in this homework that were on the second examination. I have concentrated on problems which many students did badly. They will likely be repeated in the next examination or the final examination.

$\mathcal{P}$ means $\mathcal{P}$ -TIME.

1. True/False. If the answer is not known to science at this time, enter "O" for Open.

- (a) _____ $\mathcal{P} = \mathcal{NC}$
- (b) _____ Given any $\mathcal{P}$ -TIME problem P , there is a $\mathcal{NC}$ reduction of P to the circuit value problem.
- (c) _____ Every context-free language is $\mathcal{NC}$.
- (d) _____ If a language L is generated by a general grammar, L must be decidable.
- (e) _____ Every undecidable language is $\mathcal{NP}$ -hard.

2. Draw a DFA equivalent to the following NFA.



3. Give a definition of the class $\mathcal{P}$ -complete, and name a language in that class.

4. Every language, or problem, falls into exactly one of these categories. For each of these languages or problems, write a letter indicating the correct category. You will need to search the internet for some of these.

A Known to be $\mathcal{NC}$.

B Known to be $\mathcal{P}$ -TIME, but not known to be $\mathcal{NC}$.

C Known to be $\mathcal{NP}$, but not known to be $\mathcal{P}$ -TIME and not known to be $\mathcal{NP}$ -complete.

D Known to be $\mathcal{NP}$ -complete.

E Known to be $\mathcal{P}$ -SPACE but not known to be $\mathcal{NP}$

F Known to be EXP-TIME but not known to be $\mathcal{P}$ -SPACE.

G Known to be EXP-SPACE but not known to be EXP-TIME.

H Known to be decidable, but not known to be EXP-SPACE.

I $\mathcal{RE}$ but not decidable.

K co- $\mathcal{RE}$ but not decidable.

L Neither $\mathcal{RE}$ nor co- $\mathcal{RE}$.

- (i) ----- The circuit value problem (CVP).
- (ii) ----- All C++ programs which halt with no input.
- (iii) ----- All base 10 numerals for perfect squares.
- (iv) ----- All configurations of RUSH HOUR from which it's possible to win.
- (v) ----- All satisfiable Boolean expressions.
- (vi) ----- Binary numerals for composite integers. (Composite means not prime.)
- (vii) ----- Decimal numerals for positive multiples of seven, such as 7, 14, 21, 28, ...
- (viii) ----- The furniture mover's problem. Given a room with a door and a set of furniture, is it possible to move all the furniture into the room through that door?
- (ix) ----- The set of all positions of Chinese GO, on a board of any size, from which white can win. (You will need to look this up.)
- (x) ----- The Dyck language.
- (xi) The Jigsaw problem. (That is, given a finite set of two-dimensional pieces, can they be assembled into a rectangle, with no overlap and no spaces.)
- (xii) ----- Factorization of binary numerals.
- (xiii) ----- Boolean matrix multiplication.
- (xiv) ----- All C++ programs which do not halt if given themselves as input.
- (xv) ----- SAT.
- (xvi) ----- 3-SAT.

- (xvii) ----- 2-SAT.
- (xviii) ----- The Independent Set problem.
- (xix) ----- The Subset Sum Problem.
- (xx) ----- The Block sorting problem.
- (xxi) ----- The Sliding block problem.
- (xxii) ----- The Hamiltonian cycle problem.
- (xxiii) ----- The Traveling salesman problem.
- (xxiv) ----- The Graph isomorphism problem.
- (xxv) ----- The 3-coloring problem.
- (xxvi) ----- The 2-coloring problem.
- (xxvii) ----- TOT, the set of all machine descriptions $\langle M \rangle$ such that M halts on all possible inputs.
5. Fill in the ACTION table and GOTO table of an LALR parser or the following annotated context-free grammar. Your table should enforce the convention that both operations are left-associative, and that multiplication has higher precedence than subtraction, where x represents any identifier.
1. $E \rightarrow E -_2 E_3$
 2. $E \rightarrow E *_4 E_5$
 3. $E \rightarrow x_6$

	x	$-$	$*$	$\$$	E
0					
1					
2					
3					
4					
5					
6					

6. Consider the following annotated CF grammar G .

1. $E \rightarrow E -_2 E_3$
2. $E \rightarrow -_4 E_5$
3. $E \rightarrow E \wedge_6 E_7$
4. $E \rightarrow ({}_8 E_9)_{10}$
5. $E \rightarrow x_{11}$

Here is an LALR parser for G . Walk through the computation of the parser for the input string $x - -(-x \wedge x \wedge -x)$

	x	$-$	$\wedge$	$($	$)$	$\$$	
0	s11	s4		s8			1
1		s2	s6				
2	s11	s4		s8			3
3		r1	s6		r1	r1	
4	s11	s4		s8			5
5		r2	r2		r2	r2	
6	s11	s4		s8			7
7		r3	s6		r3	r3	
8	s11	s4		s8			9
9		s2	s6		s10		
10		r4	r4		r4	r4	
11		r5	r5		r5	r5	

7. Let N be a large integer, and $w = \langle N \rangle$ its binary numeral. Let $n = |w|$. Explain how $O(n)$ parallel processors can decide, in $O(\log n)$ time, whether N is divisible by 3.

Example: in this case, N is a divisible by 3.

1011111111000011101010111100110011010100010011011111100010001011110000001011110001111110010111000010

I will give a hint for this problem in class on Thursday. Be there!

8. State the pumping lemma for regular languages. Do not get confused this time!

9. Prove that $L = \{a^n b^n : n \geq 0\}$ is not regular. Hint: By contradiction, using the pumping lemma.

10. Give a polynomial time reduction of 3-SAT to the independent set problem.

11. Prove that the halting problem is undecidable.