

## University of Nevada, Las Vegas Computer Science 456/656 Spring 2026

Additional Problems to Work in Preparation for the Second Test.

If  $w$  is any string, we write  $w^{-1}$  for the *inverse* of  $w$ . For example, if  $w = abcca$  then  $w^{-1} = accba$ . A *palindrome* is a string which is equal to its own inverse, such as “level.” We write  $\#_0(w)$  for the number of zeros in  $w$ , and  $\#_1(w)$  for the number of ones.

1. True or False. If the answer is unknown to science at this time, write **O** for Open.
  - (a) ----- Let  $L$  be the language over  $\{a, b, c\}$  consisting of all strings which have more  $a$ 's than  $b$ 's and more  $b$ 's than  $c$ 's. There is some PDA that accepts  $L$ .
  - (b) ----- The union of any two context-free languages must be context-free
  - (c) ----- The language  $\{a^n b^n \mid n \geq 0\}$  is context-free.
  - (d) ----- The language  $\{a^n b^n c^n \mid n \geq 0\}$  is context-free.
  - (e) ----- The language  $\{a^i b^j c^k \mid j = i + k\}$  is context-free.
  - (f) ----- The intersection of any regular language with any context-free language is context-free.
  - (g) ----- The intersection of any two context-free languages is context-free.
  - (h) ----- If  $L$  is a context-free language over an alphabet with just one symbol, then  $L$  is regular.
  - (i) ----- There is an LALR parser for any context-free grammar.
  - (j) ----- The set of strings that your high school algebra teacher would accept as legitimate expressions is a context-free language.
  - (k) ----- The Kleene closure of any context-free language is context-free.
  - (l) ----- Every regular language is context-free.
  - (m) ----- Every context-free language is in  $\mathcal{P}$ .
  - (n) ----- Every language accepted by a non-deterministic machine is accepted by some deterministic machine.
  - (o) ----- The problem of whether a given string is generated by a given context-free grammar is decidable.
  - (p) ----- If  $G$  is a context-free grammar, the question of whether  $L(G) = \emptyset$  is decidable.
  - (q) ----- If  $L_1$  is  $\mathcal{NP}$  complete and  $L_2$  is  $\mathcal{NP}$ , there is a  $\mathcal{P}$ -TIME reduction from  $L_1$  to  $L_2$ .
  - (r) ----- If  $L_1$  is  $\mathcal{NP}$ -hard and there is a  $\mathcal{P}$ -TIME reduction of  $L_1$  to  $L_2$ , then  $L_2$  must be  $\mathcal{NP}$ -hard.
  - (s) ----- The partition problem is  $\mathcal{NP}$  complete.

2. Give two context-free languages whose intersection is not context-free.

3. Identify which grammar generates each language. In each case, the language is over the binary alphabet  $\Sigma = \{0, 1\}$ .

(a) The language of all palindromes of even length over  $\Sigma$ .

(b) The language of all palindromes of odd length over  $\Sigma$ .

(c) The language of all strings  $w$  over  $\Sigma$  such that  $\#_0(w) = \#_1(w)$ .

(d) The language of all strings  $w$  over  $\Sigma$  such that  $\#_0(w) = \#_1(w)$  and each prefix of  $w$  has at least as many zeros as ones.

(e) The set of all binary numerals for multiples of three, where leading zeros are allowed.

(i)

$S \rightarrow 0S1S$

$S \rightarrow \lambda$

(ii)

$S \rightarrow 0S0$

$S \rightarrow 1S1$

$S \rightarrow \lambda$

(iii)

$S \rightarrow 0S0$

$S \rightarrow 1S1$

$S \rightarrow 0$

$S \rightarrow 1$

(iv)

$S \rightarrow 0S1S$

$S \rightarrow 1S0S$

$S \rightarrow \lambda$

(v)

$S \rightarrow 0S$

$S \rightarrow 1A$

$A \rightarrow 1S$

$A \rightarrow 0B$

$B \rightarrow 1B$

$B \rightarrow 0A$

$S \rightarrow 0$

4. The following grammar  $G_1$  generates a simple algebraic language. Prove that  $G_1$  is ambiguous by writing two different rightmost derivations for some string in  $L(G_1)$ . The start symbol of  $G_1$  is  $E$ .

1.  $E \rightarrow E + E$
2.  $E \rightarrow E * E$
3.  $E \rightarrow (E)$
4.  $E \rightarrow x$
5.  $E \rightarrow y$
6.  $E \rightarrow z$

5. The grammar  $G_2$ , given below, is equivalent to  $G_1$ , but is unambiguous. Draw a parse tree for  $(x + y) * z + y + x$  using  $G_2$ . ( $E$  stands for “expression,”  $T$  stands for “term,” and  $F$  stands for “factor.”)

1.  $E \rightarrow E + T$
2.  $E \rightarrow T$
3.  $T \rightarrow T * F$
4.  $T \rightarrow F$
5.  $F \rightarrow (E)$
6.  $F \rightarrow x$
7.  $F \rightarrow y$
8.  $F \rightarrow z$

6. Consider the following list of languages. In each case, the alphabet is  $\Sigma = \{a, b\}$ .

- (a) The language  $L_1$  where each prefix of any string in  $L_2$  has at least as many  $a$ 's as  $b$ 's.
- (b) The language  $L_2 \subseteq L_1$  consisting of all strings which have an equal number of each symbol.
- (c) The language of all even length palindromes.

Each of the above languages is accepted by one of the PDA shown below. Which one?

