

True/False Questions

\mathcal{P} means \mathcal{P} -TIME

\mathcal{NP} means \mathcal{NP} -TIME

\mathcal{RE} means recursively enumerable

\mathcal{NC} means Nick's class.

If \mathcal{C} is any class of languages, $\text{co-}\mathcal{C}$ means the class of all languages which are complements of languages in \mathcal{C} .

A *binary language* is a language over the binary alphabet $\{0, 1\}$.

A *recursive* function is any function which can be computed by a machine.

A *recursive* real number is any real number whose n^{th} decimal digit is a recursive function of n .

A *fraction* is a string consisting of a numeral, followed by a slash, followed by another numeral.

1. True or False. T = true, F = false, and O = open, meaning that the answer is not known science at this time. In the questions below, \mathcal{P} and \mathcal{NP} denote \mathcal{P} -TIME and \mathcal{NP} -TIME, respectively.
 - (i) ----- The problem of whether a given string is generated by a given context-free grammar is decidable.
 - (ii) ----- If G is a context-free grammar, the question of whether $L(G) = \emptyset$ is decidable.
 - (iii) ----- The Kleene closure of any \mathcal{NP} language is \mathcal{NP}
 - (iv) ----- The language $\{a^n b^n c^n d^n \mid n \geq 0\}$ is recursive.
 - (v) ----- The language $\{a^n b^n c^n \mid n \geq 0\}$ is in the class \mathcal{P} -TIME.
 - (vi) ----- There exists a polynomial time algorithm which finds the factors of any positive integer, where the input is given as a binary numeral.
 - (vii) ----- Every undecidable problem is \mathcal{NP} -complete.
 - (viii) ----- Every problem that can be mathematically defined has an algorithmic solution.
 - (ix) ----- The intersection of two undecidable languages is always undecidable.
 - (x) ----- Every \mathcal{NP} language is decidable.
 - (xi) ----- If L_1 and L_2 are \mathcal{NP} -complete languages and $L_1 \cap L_2$ is not empty, then $L_1 \cap L_2$ must be \mathcal{NP} -complete.
 - (xii) ----- There exists a \mathcal{P} -TIME algorithm which finds a maximum independent set in any graph G .
 - (xiii) ----- There exists a \mathcal{P} -TIME algorithm which finds a maximum independent set in any acyclic graph G .
 - (xiv) ----- $\mathcal{NC} = \mathcal{P}$.
 - (xv) ----- $\mathcal{P} = \mathcal{NP}$.
 - (xvi) ----- $\mathcal{NP} = \mathcal{P}$ -SPACE

- (xvii) ----- \mathcal{P} -SPACE = EXP-TIME
- (xviii) ----- EXP-TIME = EXP-SPACE
- (xix) ----- The traveling salesman problem (TSP) is known to be \mathcal{NP} -complete.
- (xx) ----- The language consisting of all satisfiable Boolean expressions is known to be \mathcal{NP} -complete.
- (xxi) ----- The Boolean Circuit Problem is in \mathcal{P} .
- (xxii) ----- The Boolean Circuit Problem is in \mathcal{NC} .
- (xxiii) ----- 2-SAT is \mathcal{P} -TIME.
- (xxiv) ----- 3-SAT is \mathcal{P} -TIME.
- (xxv) ----- Primality is \mathcal{P} -TIME.
- (xxvi) ----- There is a \mathcal{P} -TIME reduction of the halting problem to 3-SAT.
- (xxvii) ----- Every context-free language is in \mathcal{NC} .
- (xxviii) ----- Addition of binary numerals is in \mathcal{NC} .
- (xxix) ----- Every language generated by a general grammar is recursive.
- (xxx) ----- The problem of whether two given context-free grammars generate the same language is decidable.
- (xxxi) -----
The language of all fractions (using base 10 numeration) whose values are less than π is decidable.
- (xxxii) ----- For any two languages L_1 and L_2 , if L_1 is undecidable and there is a recursive reduction of L_1 to L_2 , then L_2 must be undecidable.
- (xxxiii) ----- For any two languages L_1 and L_2 , if L_2 is undecidable and there is a recursive reduction of L_1 to L_2 , then L_1 must be undecidable.
- (xxxiv) ----- If P is a mathematical proposition that can be written using a string of length n , and P has a proof, then P must have a proof whose length is $O(2^{2^n})$.
- (xxxv) ----- If L is any \mathcal{NP} language, there must be a \mathcal{P} -TIME reduction of L to the partition problem.
- (xxxvi) ----- If L is \mathcal{NP} and also co- \mathcal{NP} , then L must be \mathcal{P} .
- (xxxvii) ----- A language is \mathcal{RE} if and only if it is generated by a grammar.
- (xxxviii) ----- If L is \mathcal{RE} and also co- \mathcal{RE} , then L must be decidable.
- (xxxix) ----- Every language is enumerable.
- (xl) ----- If a language L is undecidable, then there can be no machine that enumerates L .

- (xli) ----- There exists a mathematical proposition which is true, but can be neither proved nor disproved.
 - (xlii) ----- There is a non-recursive function which grows faster than any recursive function.
 - (xliii) ----- There exists a machine that runs forever and outputs the string of decimal digits of π (the well-known ratio of the circumference of a circle to its diameter).
 - (xliv) ----- For every real number x , there exists a machine that runs forever and outputs the string of decimal digits of x .
 - (xlv) ----- Every subset of any enumerable set is enumerable.
 - (xlvi) ----- There is a polynomial time reduction of the subset sum problem to the binary numeral factorization problem.
 - (xlvii) ----- For any real number x , the set of fractions whose values are less than x is \mathcal{RE} .
 - (xlviii) ----- For any recursive real number x , the set of fractions whose values are less than x is recursive (i.e., decidable).
 - (xlix) ----- The membership problem for any CFL is in the class \mathcal{P} -TIME.
 - (l) ----- 2SAT is known to be \mathcal{NP} -complete.
 - (li) ----- The complement of any \mathcal{P} -TIME language is \mathcal{P} -TIME.
 - (lii) ----- The complement of any \mathcal{P} -SPACE language is \mathcal{P} -SPACE.
 - (liii) ----- The complement of any decidable language is decidable.
 - (liv) ----- The complement of any undecidable language is undecidable.
 - (lv) ----- The complement of any \mathcal{RE} language is \mathcal{RE} .
- The *jigsaw puzzle problem* is, given a set of various polygons, and given a rectangular table, is it possible to assemble those polygons to exactly cover the table?
- The *furniture mover's problem* is, given a room with a door, and given a set of objects outside the room, is it possible to move all the objects into the room through the door?
- (lvi) ----- The jigsaw puzzle problem is known to be \mathcal{NP} complete.
 - (lvii) ----- The jigsaw puzzle problem is known to be \mathcal{P} -SPACE complete.
 - (lviii) ----- The furniture mover's problem is known to be \mathcal{NP} complete.
 - (lix) ----- The furniture mover's problem is known to be \mathcal{P} -SPACE complete.
 - (lx) ----- The complement of any recursive language is recursive.
 - (lxi) ----- For any infinite countable sets A and B , there is a 1-1 correspondence between A and B .

- (lxii) ----- The set of all binary languages is countable.
- (lxiii) ----- A language L is recursively enumerable if and only if there is a machine which accepts L .
- (lxiv) ----- Every \mathcal{NP} language is reducible to the independent set problem in polynomial time.
- (lxv) ----- If a Boolean expression is satisfiable, there is a polynomial time proof that it is satisfiable.
- (lxvi) ----- The general sliding block problem is \mathcal{P} -SPACE complete.
- (lxvii) ----- The halting problem is decidable.