## True/False Questions

\$\mathcal{P}\$ means \$\mathcal{P}\$-TIME
\$\mathcal{N}\mathcal{P}\$-TIME
\$\mathcal{R}\mathcal{E}\$ means recursively enumerable
\$\mathcal{N}\mathcal{C}\$ means Nick's class.
If \$\mathcal{C}\$ is any class of languages, co-\$\mathcal{C}\$ means the class of all languages which are complements of languages in \$\mathcal{C}\$.
A binary language is a language over the binary alphabet {0,1}.
A recursive function is any function which can be computed by a machine.
A recursive real number is any real number whose \$n^{th}\$ decimal digit is a recursive function of \$n\$.

- A *fraction* is a string consisting of a numeral, followed by a slash, followed by another numeral.
  - 1. True or False. T = true, F = false, and O = open, meaning that the answer is not known science at this time. In the questions below,  $\mathcal{P}$  and  $\mathcal{NP}$  denote  $\mathcal{P}$ -TIME and  $\mathcal{NP}$ -TIME, respectively.
    - (i) \_\_\_\_\_ The problem of whether a given string is generated by a given context-free grammar is decidable.
    - (ii) \_\_\_\_\_ If G is a context-free grammar, the question of whether  $L(G) = \emptyset$  is decidable.
    - (iii) \_\_\_\_\_ The Kleene closure of any  $\mathcal{NP}$  langauge is  $\mathcal{NP}$
    - (iv) \_\_\_\_\_ The language  $\{a^n b^n c^n d^n \mid n \ge 0\}$  is recursive.
    - (v) \_\_\_\_\_ The language  $\{a^n b^n c^n \mid n \ge 0\}$  is in the class  $\mathcal{P}$ -TIME.
    - (vi) \_\_\_\_\_ There exists a polynomial time algorithm which finds the factors of any positive integer, where the input is given as a binary numeral.
    - (vii)  $\_$  Every undecidable problem is  $\mathcal{NP}$ -complete.
    - (viii) \_\_\_\_\_ Every problem that can be mathematically defined has an algorithmic solution.
    - (ix) \_\_\_\_\_ The intersection of two undecidable languages is always undecidable.
    - (x) \_\_\_\_\_ Every  $\mathcal{NP}$  language is decidable.
    - (xi) \_\_\_\_\_ If  $L_1$  and  $L_2$  are  $\mathcal{NP}$ -complete languages and  $L_1 \cap L_2$  is not empty, then  $L_1 \cap L_2$  must be  $\mathcal{NP}$ -complete.
    - (xii)  $\ldots$  There exists a  $\mathcal{P}$ -TIME algorithm which finds a maximum independent set in any graph G.
    - (xiii) \_\_\_\_\_ There exists a  $\mathcal{P}$ -TIME algorithm which finds a maximum independent set in any acyclic graph G.
    - (xiv)  $\ldots \mathcal{NC} = \mathcal{P}.$
    - (xv)  $\mathcal{P} = \mathcal{N}\mathcal{P}.$
    - (xvi)  $\dots \mathcal{NP} = \mathcal{P}\text{-space}$

- (xvii)  $\dots \mathcal{P}$ -Space = EXP-time
- (xviii) \_\_\_\_\_ EXP-TIME = EXP-SPACE
- (xix)  $\_$  The traveling salesman problem (TSP) is known to be  $\mathcal{NP}$ -complete.
- (xx)  $\dots$  The language consisting of all satisfiable Boolean expressions is known to be  $\mathcal{NP}$ -complete.
- (xxi)  $\_$  The Boolean Circuit Problem is in  $\mathcal{P}$ .
- (xxii)  $\_$  The Boolean Circuit Problem is in  $\mathcal{NC}$ .
- (xxiii)  $\_\_\_2$ -SAT is  $\mathcal{P}$ -TIME.
- (xxiv)  $\_\_\_$  3-SAT is  $\mathcal{P}$ -TIME.
- (xxv) ------ Primality is  $\mathcal{P}$ -TIME.
- (xxvi)  $\_$  There is a  $\mathcal{P}$ -TIME reduction of the halting problem to 3-SAT.
- (xxvii)  $\_$  Every context-free language is in  $\mathcal{NC}$ .
- (xxviii)  $\_$  Addition of binary numerals is in  $\mathcal{NC}$ .
- (xxix) \_\_\_\_\_ Every language generated by a general grammar is recursive.
- (xxx) \_\_\_\_\_ The problem of whether two given context-free grammars generate the same language is decidable.
- (xxxi) \_\_\_\_\_ The language of all fractions (using base 10 numeration) whose values are less than  $\pi$  is decidable.
- (xxxii) \_\_\_\_\_\_ For any two languages  $L_1$  and  $L_2$ , if  $L_1$  is undecidable and there is a recursive reduction of  $L_1$  to  $L_2$ , then  $L_2$  must be undecidable.
- (xxxiii) \_\_\_\_\_\_ For any two languages  $L_1$  and  $L_2$ , if  $L_2$  is undecidable and there is a recursive reduction of  $L_1$  to  $L_2$ , then  $L_1$  must be undecidable.
- (xxxiv) \_\_\_\_\_ If P is a mathematical proposition that can be written using a string of length n, and P has a proof, then P must have a proof whose length is  $O(2^{2^n})$ .
- (xxxv) \_\_\_\_\_ If L is any  $\mathcal{NP}$  language, there must be a  $\mathcal{P}$ -TIME reduction of L to the partition problem.
- (xxxvi) \_\_\_\_\_ If L is  $\mathcal{NP}$  and also co- $\mathcal{NP}$ , then L must be  $\mathcal{P}$ .
- (xxxvii) \_\_\_\_\_ A language is  $\mathcal{RE}$  if and only if it is generated by a grammar.
- (xxxviii) \_\_\_\_\_ If L is  $\mathcal{RE}$  and also co- $\mathcal{RE}$ , then L must be decidable.
- (xxxix) \_\_\_\_\_ Every language is enumerable.
  - (xl)  $\ldots$  If a language L is undecidable, then there can be no machine that enumerates L.

- (xli) \_\_\_\_\_ There exists a mathematical proposition which is true, but can be neither proved nor disproved.
- (xlii) \_\_\_\_\_ There is a non-recursive function which grows faster than any recursive function.
- (xliii) \_\_\_\_\_ There exists a machine that runs forever and outputs the string of decimal digits of  $\pi$  (the well-known ratio of the circumference of a circle to its diameter).
- (xliv) \_\_\_\_\_ For every real number x, there exists a machine that runs forever and outputs the string of decimal digits of x.
- (xlv) \_\_\_\_\_ Every subset of any enumerable set is enumerable.
- (xlvi) \_\_\_\_\_ There is a polynomial time reduction of the subset sum problem to the binary numeral factorization problem.
- (xlvii) \_\_\_\_\_ For any real number x, the set of fractions whose values are less than x is  $\mathcal{RE}$ .
- (xlviii) \_\_\_\_\_ For any recursive real number x, the set of fractions whose values are less than x is recursive (i.e., decidable).
- (xlix) \_\_\_\_\_ The membership problem for any CFL is in the class  $\mathcal{P}$ -TIME.
  - (1) \_\_\_\_\_ 2-SAT is known to be  $\mathcal{NP}$ -complete.
  - (li) \_\_\_\_\_ The complement of any  $\mathcal{P}$ -TIME language is  $\mathcal{P}$ -TIME.
  - (lii) \_\_\_\_\_ The complement of any  $\mathcal{P}$ -SPACE language is  $\mathcal{P}$ -SPACE.
- (liii) \_\_\_\_\_ The complement of any decidable language is decidable.
- (liv) \_\_\_\_\_ The complement of any undecidable language is undecidable.
- (lv)  $\ldots$  The complement of any  $\mathcal{RE}$  language is  $\mathcal{RE}$ .

The *jigsaw puzzle problem* is, given a set of various polygons, and given a rectangular table, is it possible to assemble those polygons to exactly cover the table?

The *furniture mover's problem* is, given a room with a door, and given a set of objects outside the room, it is possible to move all the objects into the room through the door?

- (lvi) \_\_\_\_\_ The jigsaw puzzle problem is known to be  $\mathcal{NP}$  complete.
- (lvii)  $\_$  The jigsaw puzzle problem is known to be  $\mathcal{P}$ -SPACE complete.
- (lviii) \_\_\_\_\_ The furniture mover's problem is known to be  $\mathcal{NP}$  complete.
- (lix) \_\_\_\_\_ The furniture mover's problem is known to be  $\mathcal{P}$ -SPACE complete.
- (lx) \_\_\_\_\_ The complement of any recursive language is recursive.
- (lxi) \_\_\_\_\_ For any infinite countable sets A and B, there is a 1-1 correspondence between A and B.

- (lxii) \_\_\_\_\_ The set of all binary languages is countable.
- (lxiii)  $\_\_\_\_\_$  A language L is recursively enumerable if and only if there is a machine which accepts L.
- (lxiv) \_\_\_\_\_ Every  $\mathcal{NP}$  language is reducible to the independent set problem in polynomial time.
- (lxv) \_\_\_\_\_ If a Boolean expression is satisfiable, there is a polynomial time proof that it is satisfiable.
- (lxvi)  $\_$  The general sliding block problem is  $\mathcal{P}$ -SPACE complete.
- (lxvii) \_\_\_\_\_The halting problem is decidable.